

Test and Verification Lecture 14

SPIN and promela

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Plan for today

- Promela
 - Constructs
 - Examples
- LTL properties
- Installation
- SPIN demo

Promela

- Programming Meta Language
- A modeling language for verification and simulation
- Restricted set of constructs and datatypes

Model parts

- Processes
- Message Channels
- Variables

Executability

- No difference between conditions and statements
 - This might seem strange at first
- Boolean conditions can be executed when they are true
- Else they block until they become true
- Statement are always executable

Executability

- No need for busy loops

```
while (a != b)
    skip    /* wait for a==b */
```

- Can be replaced with

```
(a == b)
```

Variables

- Global and local variables

```
bool flag;  
int state;  
byte msg;
```

- Array variables
- Message types

```
mtype = {ack, nack, err}
```

Datatypes

Type name	C-equivalent	Macro in limits.h	Typical Range
bit or bool	bit-field	-	0 .. 1
byte	uchar	CHAR_BIT (width in bits)	0 .. 255
short	short	SHRT_MIN..SHRT_MAX	$-2^{15} - 1 \dots 2^{15} - 1$
int	int	INT_MIN..INT_MAX	$-2^{31} - 1 \dots 2^{31} - 1$

Proctype

- One local variable

```
proctype A( )  
{  
    byte state;  
    state = 3  
}
```

Proctype

- ; is only a separator
- -> is equivalent

```
byte state = 2;  
proctype A()  
{  
    (state == 1) -> state = 3  
}  
proctype B()  
{  
    state = state - 1  
}
```

Process Instantiation

- Special init process

```
init
{
    run A();
    run B()
}
```

- Processes can be started from anywhere

Passing variables

```
proctype A(byte state; short foo)
{
    (state == 1) -> state = foo
}
init
{
    run A(1, 3)
}
```

Mutual exclusion example

```
#define true    1
#define false   0
#define Aturn   false
#define Bturn   true
bool x, y, t;
proctype A()
{
    x = true;
    t = Bturn;
    (y == false || t == Aturn);
    /* critical section */
    x = false
}
proctype B()
{
    y = true;
    t = Aturn;
    (x == false || t == Bturn);
    /* critical section */
    y = false
}
init
{
    run A(); run B()
}
```

Atomic sequences

- Runtime error if anything but the first statement blocks

```
byte state = 1;
proctype A()
{
    atomic {
        (state==1) -> state = state+1
    }
}
proctype B()
{
    atomic {
        (state==1) -> state = state-1
    }
}
init
{
    run A(); run B()
}
```

Message passing

- Used to model transfer of data
- Global or local
- Channels can send channel names

```
chan qname = [16] of { short }
chan qname = [16] of { byte, int, chan, byte }
```

- Synchronous communication

```
chan qname = [0] of { short }
```

Message passing

- **Sending**

```
chan qname = [16] of { byte, int, chan, byte }
```

```
qname!v,y,myChan,a
```

- **Receiving**

```
qname?var,x,ch,b
```

- **Receiving with constants**

```
qname?var,cons1,ch,cons2
```

Example

```
proctype A(chan q1)
{
    chan q2;
    q1?q2;
    q2!123
}
proctype B(chan qforb)
{
    int x;
    qforb?x;
    printf( "x = %d\n" , x )
}
init {
    chan qname = [1] of { chan };
    chan qforb = [1] of { int };
    run A(qname);
    run B(qforb);
    qname!qforb
}
```

Testing for messages

- Length – built in function

```
len(qname)
```

- Testing for reception

```
qname? [var, cons1, ch, cons2]
```

- True if the message can be received
- Remember to use atomic

```
(len(qname) < MAX) -> qname!msgtype  
qname? [msgtype] -> qname?msgtype
```

Control Flow

- Case selection

```
if
::: (a != b) -> option1
::: (a == b) -> option2
fi
```

- Guards
- Does not need to be mutually exclusive
- Keyword else

Repetition

```
proctype counter()
{
    do
        :: (count != 0) ->
            if
                :: count = count + 1
                :: count = count - 1
            fi
        :: (count == 0) -> break
    od
}
```

Unconditional Jumps

```
proctype Euclid(int x, y)
{
    do
        :: (x > y) -> x = x - y
        :: (x < y) -> y = y - x
        :: (x == y) -> goto done
    od;
done:
    skip
}
```

- Extra skip at the end

Return values

```
proctype fact(int n; chan p)
{
    chan child = [1] of { int };
    int result;
    if
        :: (n <= 1) -> p!1
        :: (n >= 2) ->
            run fact(n-1, child);
            child?result;
            p!n*result
    fi
}
init
{
    chan child = [1] of { int };
    int result;
    run fact(7, child);
    child?result;
    printf("result: %d\n", result)
}
```

Timeout

- Modeling trick

```
proctype watchdog( )  
{  
    do  
        :: timeout -> guard!reset  
    od  
}
```

- Cannot be implemented

Assertions

- Produces errors during simulation or verification

```
assert(any_boolean_condition)
```

Labels

- End state labels
 - end, end1, end_here, ...
- Progress
 - progress, progress2, ...
- After having compiled
 - ./pan -l
 - Search for non progress loops

SPIN

- spin -m -a ex.1a
- gcc -o pan pan.c
- ./pan

Bitstate hashing

- Coverage
- Not precise analysis
- -DBITSTATE

- Propositional formulas defined separately
 - Evaluated over computations
- [] Always
- <> Eventually
- U (strong) until ($p \wedge q$)
- V !($\neg p \wedge \neg q$) (Also known as release)

Examples

- Nested properties
- $\{\!\! \{ p$
- $!(\ <\!> !q\)$
- $p \cup q$
- $p \cup (\{\!\! \{ (q \cup r))$