

Finding Fastest Paths on A Road Network with Speed Patterns

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Outline

Background

Fastest-Path Computation

Experimental Results

Strong and Weak Points

Problem

Cause

Current navigation system assume that the travel time on a road segment is constant.

Effect

The travel-time in for example rush hour is highly underestimated.

Solution

“I may leave for work any time between 7am and 9am; please suggest all fastest paths, e.g. take route A if the leaving time is between 7 and 7:45, and take route B otherwise.”

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CapeCod Network

Definition

Day Category D is a day category set.

Example

$$D = \{workday, non - workday\}$$

Definition

Categorized Piecewise Constant speed (CapeCod) CapeCod is a pattern of one daily speed pattern for every day-category in D .

Example

$$\begin{aligned} \text{workday} = & [00 : 00 - 07 : 00) = 1.00 \text{mp/h} \\ & [07 : 00 - 09 : 00) = 0.75 \text{mp/h} \\ & [09 : 00 - 15 : 00) = 1.00 \text{mp/h} \\ & [15 : 00 - 17 : 00) = 0.60 \text{mp/h} \\ & [17 : 00 - 24 : 00) = 1.00 \text{mp/h} \end{aligned}$$

CapeCod Network, cont.

Definition

CapeCod Network A directed graph $G(N, E)$ where $N = \{(n_i, loc_i) \mid i \in [1, m]\}$ is the set of nodes and $E = \{n_i, n_j, d_{ij}, pat_{ij} \mid i, j \in [1, m]\}$ is the set of edges where d_{ij} is a distance and pat_{ij} is a CapeCod pattern.

Queries

Single Fastest Path (singleFP) Query

Given a start node s , an end node e , and a leaving time interval I , find the time instant $t \in I$ and the corresponding fastest path from s to e such that leaving from s at time t minimizes the travel time from s to e

All Fastest Path (allFP) Query

Given a CapeCod network, a start node s and an end node e , and a leaving time interval I , find a full partitioning of I : I_1, \dots, I_k , where each sub-interval is associated with a fastest path, such that two leaving time instants in one sub-interval leads to the same fastest-path and two leaving time instants in two adjacent sub-intervals leads to different fastest paths.

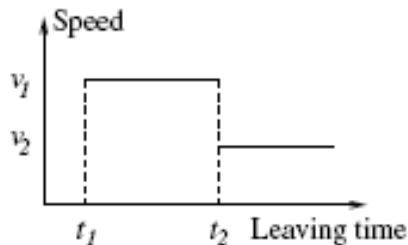
Storage Model

- ▶ Connectivity Cluster Access Method (CCAM)
- ▶ For each node n_i
 - ▶ Stores the location loc_i
 - ▶ Stores the list of neighbors
 - ▶ Stores the Euclidian distance to each neighbor
 - ▶ Stores the CapeCod speed pattern (pat_{ij}) to each neighbor
- ▶ Hilbert curve value indexes the loc_i value

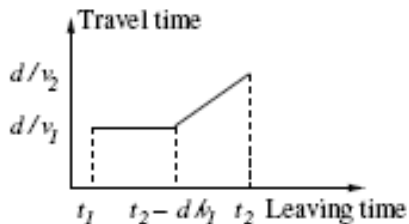
Operations

- ▶ $FindNode(n_i)$
- ▶ $GetSuccessor(n_i)$
- ▶ Modify the network

From Speed Pattern to Travel Time



(a) Speed



(b) Travel time

- ▶ Speed v_1 in $[t_1, t_2)$ and v_2 there after
- ▶ Speed can be changed more than once (not likely)

Travel-Time Function

$$T(l \in [t_1, t_2], n_i \rightarrow n_j) = \begin{cases} \frac{d}{v_1} & l \in [t_1, t_2 - \frac{d}{v_1}] \\ (1 - \frac{v_1}{v_2})(t_2 - l) + \frac{d}{v_2} & l \in [t_2 - \frac{d}{v_1}, t_2] \end{cases}$$

Outline

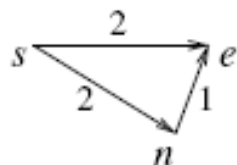
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Initialization



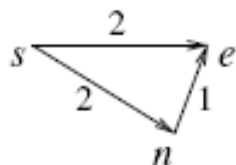
$$s \rightarrow e: [6-8):1/3$$

$$s \rightarrow n: [6-7):1/3, [7-8):1$$

$$n \rightarrow e: [6-7:08):1/3, [7:08-8):1/10$$

$$I = [6 : 50 - 7 : 05]$$

Initialization



$$s \rightarrow e: [6-8):1/3$$

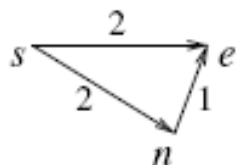
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$$T(I \in [6 : 50 - 7 : 05], s \rightarrow e) =$$

Initialization



$$s \rightarrow e: [6-8):1/3$$

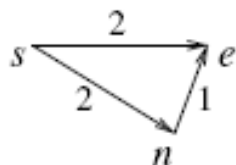
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$$I = [6 : 50 - 7 : 05]$$

$$T(I \in [6 : 50 - 7 : 05], s \rightarrow e) = 6$$

Initialization



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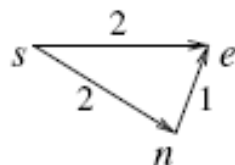
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Initialization



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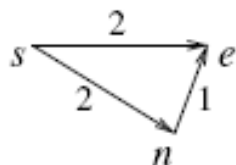
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$$I = [6 : 50 - 7 : 05]$$

$$T(I \in [6 : 50 - 7 : 05], s \rightarrow e) = 6$$

$$T(I \in [6 : 50 - 7 : 05], s \rightarrow n) = \begin{cases} 6 & I \in [6 : 50 - 6 : 54) \end{cases}$$

Initialization



$$s \rightarrow e: [6-8):1/3$$

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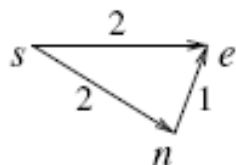
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$$I = [6 : 50 - 7 : 05]$$

$$T(I \in [6 : 50 - 7 : 05], s \rightarrow e) = 6$$

$$T(I \in [6 : 50 - 7 : 05], s \rightarrow n) = \begin{cases} 6 & I \in [6 : 50 - 6 : 54) \\ 2 & I \in [7 : 00 - 7 : 05) \end{cases}$$

Initialization



$$s \rightarrow e: [6-8):1/3$$

$$s \rightarrow n: [6-7):1/3, [7-8):1$$

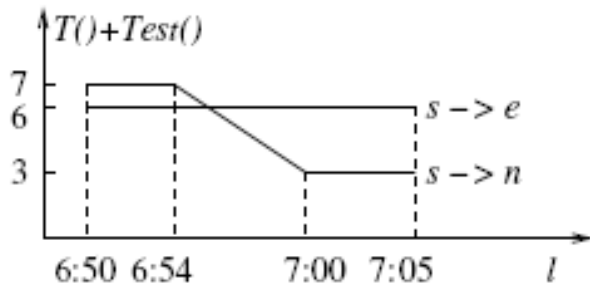
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$$I = [6 : 50 - 7 : 05]$$

$$T(I \in [6 : 50 - 7 : 05], s \rightarrow e) = 6$$

$$T(I \in [6 : 50 - 7 : 05], s \rightarrow n) = \begin{cases} 6 & I \in [6 : 50 - 6 : 54) \\ 2 & I \in [7 : 00 - 7 : 05) \\ \frac{2}{3}(7 : 00 - I) + 2 & I \in [6 : 54 - 7 : 00) \end{cases}$$

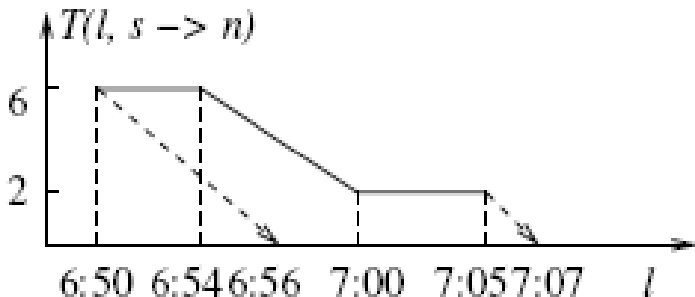
Path Expansion



Based on $T() + T_{est}()$

- ▶ Must expand the path $s \rightarrow n$

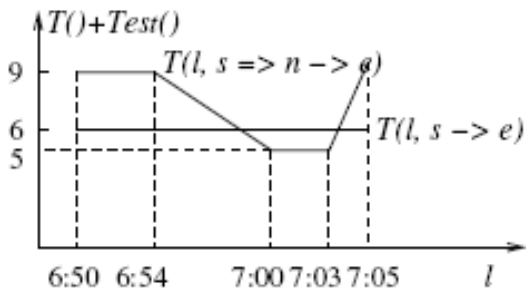
Path Expansion, cont.



- ▶ Determine the interval during which the travel-time function for road segment is needed
- ▶ Determine the time instants at which the resulting function changes
- ▶ For each time interval determine the travel-time function

$$T(l \in [6:56-7:07], n \rightarrow e) = \begin{cases} 3 & l \in [6:56 - 7:05] \\ 10 - \frac{7}{3}(7:08 - l) & l \in [7:05 - 7:07] \end{cases}$$

Fastest Path Result



$$T(l \in [6:50-7:05], s \rightarrow e) = \begin{cases} s \rightarrow e & l \in [6:50 - 6:58:30) \\ s \rightarrow n \rightarrow e & l \in [6:58:30 - 7:03:26) \\ s \rightarrow e & l \in [7:03:26 - 7:05) \end{cases}$$

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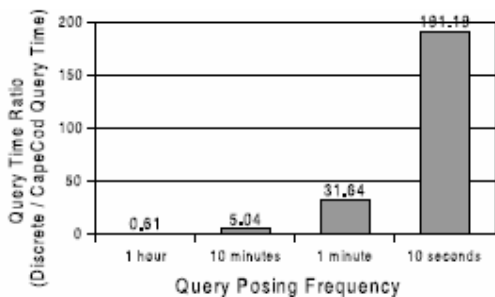
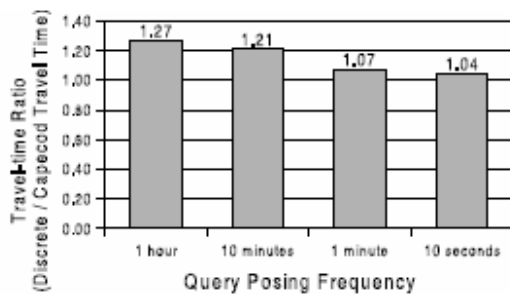
Strong and Weak Points

Setup

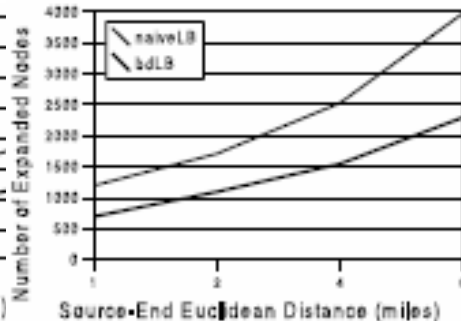
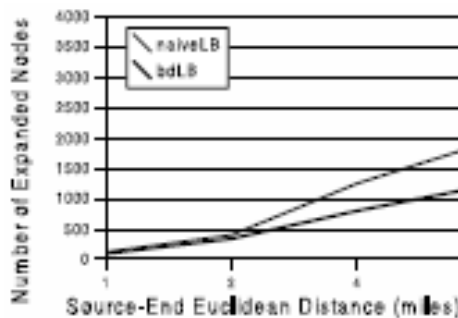
- ▶ Road network : Suffolk county of Massachusetts (Boston)
 - ▶ 14,456 nodes and 20,461 directed edges
- ▶ Day categories D : = { workday and non-workday }
- ▶ Road segments types
 - ▶ inbound highways
 - ▶ outbound highway
 - ▶ local roads outside Boston
 - ▶ local roads inside Boston
- ▶ Speed Pattern

	Inbound Hw		Outbound Hw		Local (in)		Local (out)
non-workday	65		65		40		40
workday	00:00-07:00	65	00:00-04:00	65	00:00-04:00	40	40
	07:00-10:00	20	04:00-07:00	30	04:00-07:00	20	
	10:00-24:00	65	07:00-24:00	65	07:00-24:00	40	

Performance



Performance, cont.



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Strong Points

- ▶ Contributions of paper clearly pointed out
- ▶ Good coverage of related work
- ▶ Explanation of core algorithm in Section 4 including examples

Weak Points

- ▶ Synthetic setup of CapeCod travel-time speed patterns
- ▶ Some references are pretty old ([5, 9, 10, 12])
- ▶ No validation of computed travel times to actual travel times