

UPPAAL TRON: Testing Real-time systems ONline

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Center for Embedded Software Systems



Basic Research in Computer Science



Aalborg University

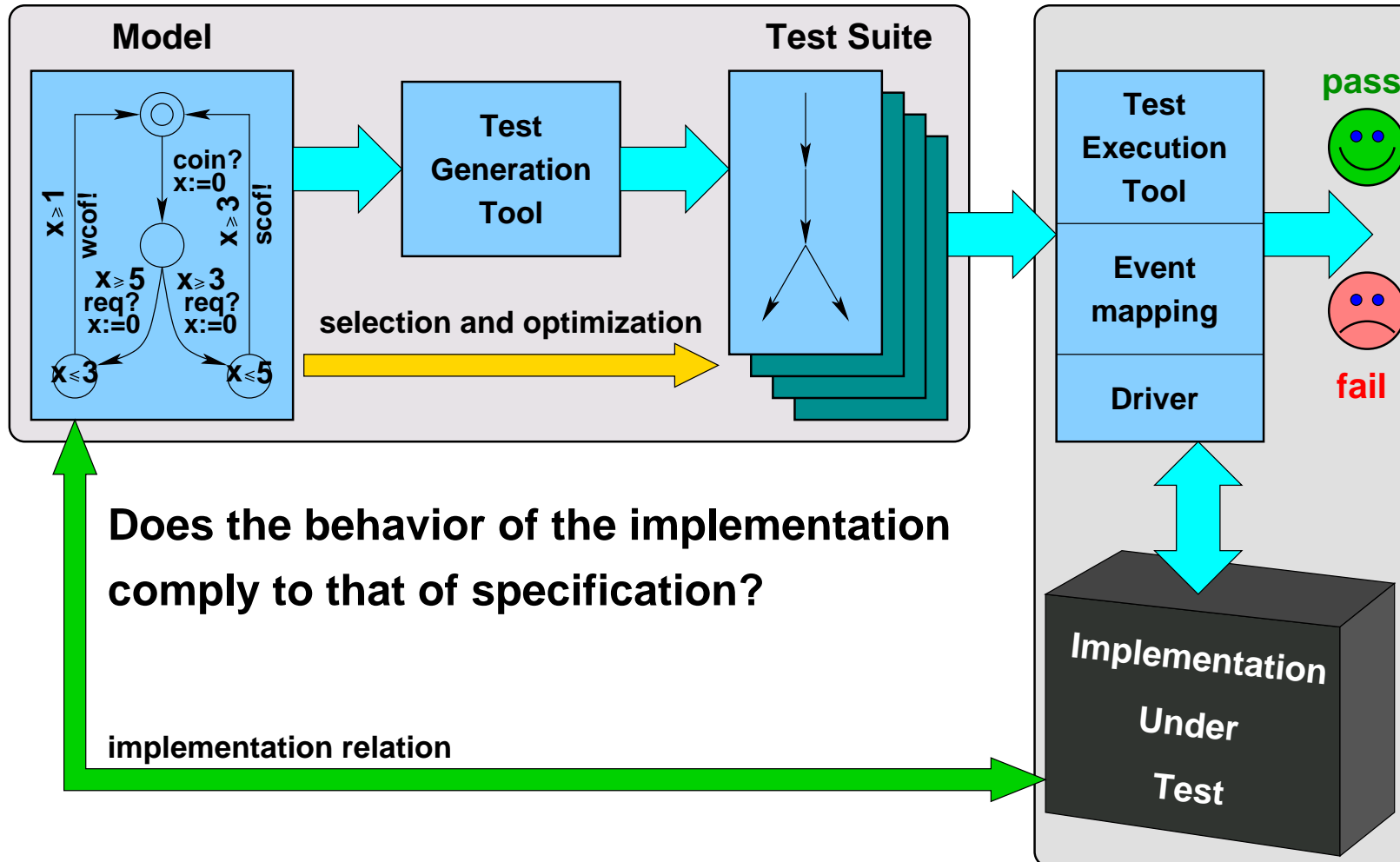
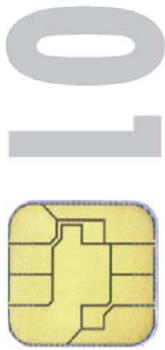


- Classical model-based testing.
- Online test setup: from specification to testing.
- UPPAAL TRON framework.
- Conformance relations: tioco and rtioco.
- Online test algorithm.
- Conclusions.
- Future work.
- Demo of Java “light controller” online test.

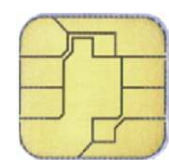
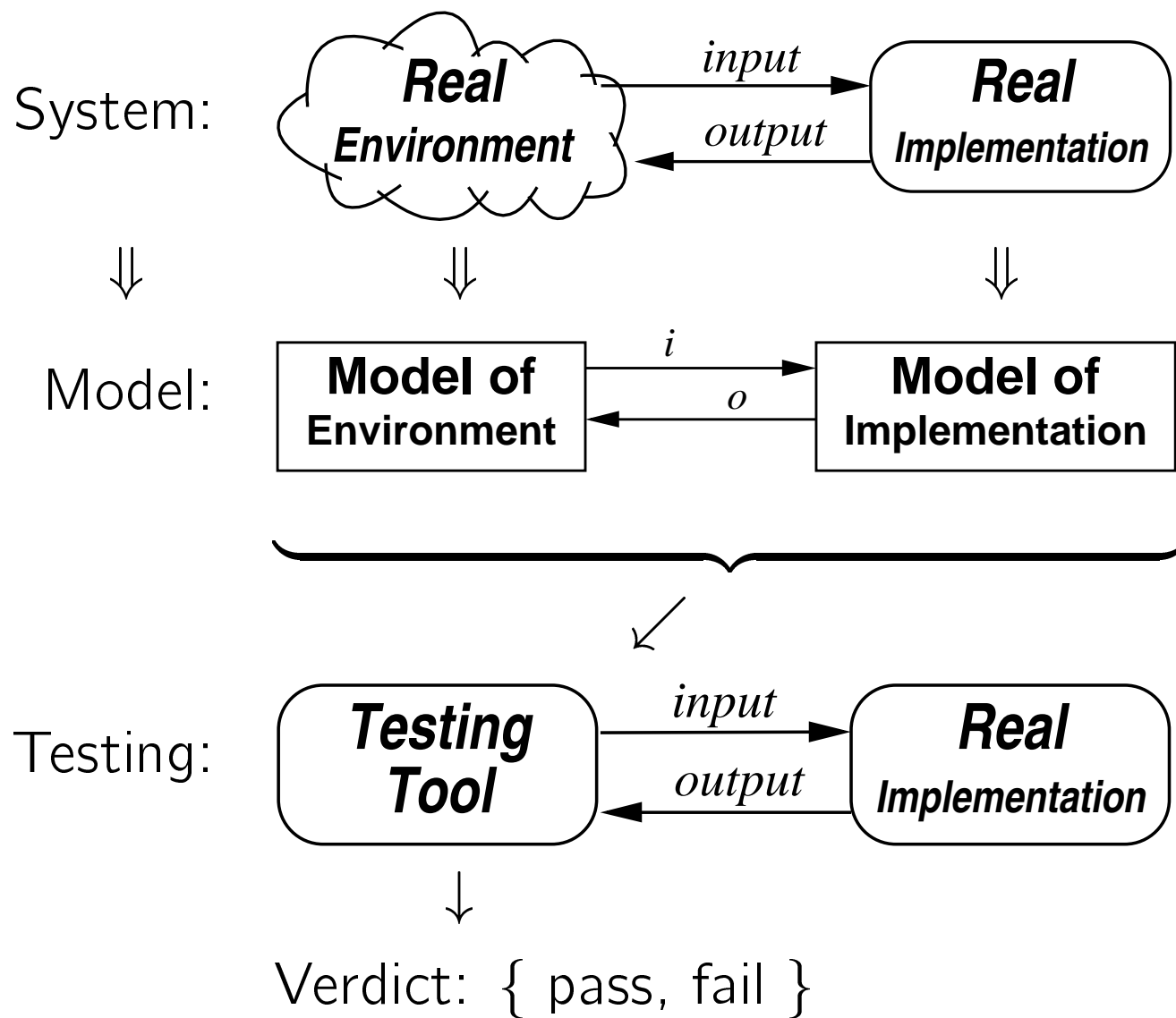


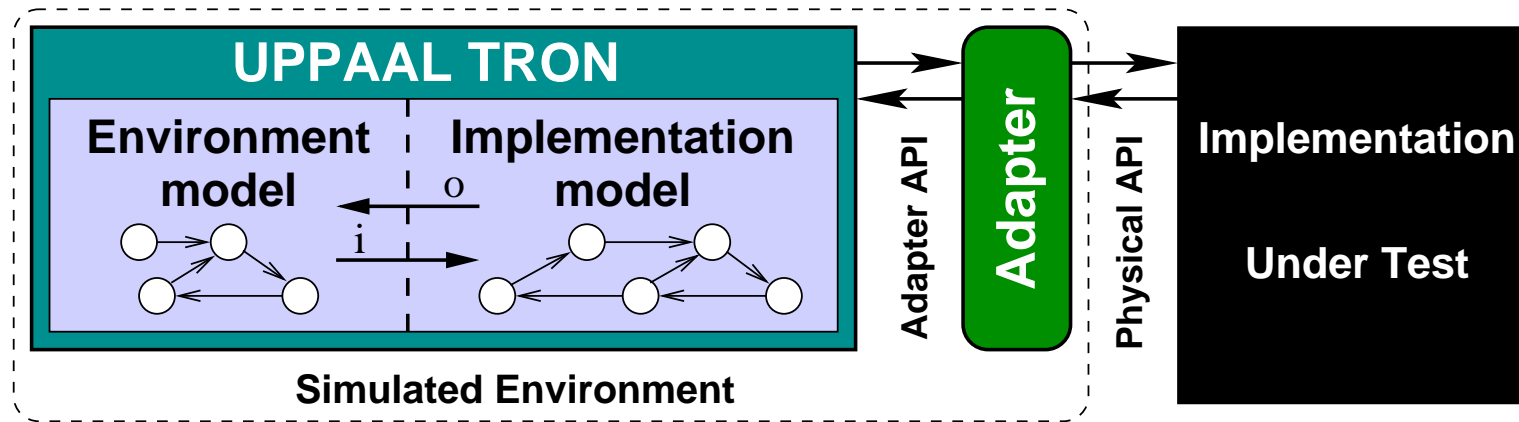
Classical Model-based Testing Framework

- Black-box, model-based, conformance testing offline.

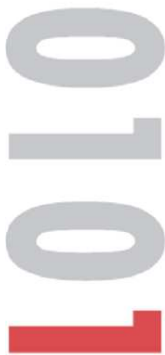
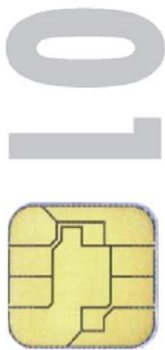


Online setup: System \Rightarrow Model \Rightarrow Testing





- UPPAAL timed automata specification:
 - Model of (requirements for) implementation,
 - Model of (assumptions about) environment,
 - Contain several concurrent non-deterministic processes.
- Online testing tool:
 - generates test primitives for stimuli (input) with timing,
 - verifies the validity of response (output) with timing,
 - while test primitives are executed online.



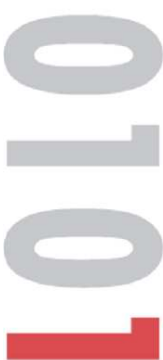
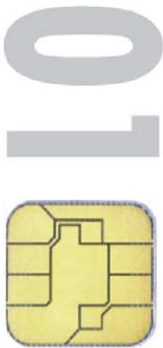
- Defines the correctness criteria by analyzing observable traces.
- Intuitively:
 - timed trace e.g.: $\sigma = coin? \cdot 5 \cdot req? \cdot 2 \cdot weakCoffee! \cdot 9 \cdot coin?$
 - observed output should be matched and allowed in specification,
 - timing of observed output should be allowed in specification,
 - any observation outside the specification is a fault.
- Formally:

- *Timed traces* from state s : $TTr(s) \stackrel{def}{=} \{\sigma \in (A \cup \mathbb{R}_{\geq 0})^* \mid s \xrightarrow{\sigma}\}$

Outputs: $Out(S) \stackrel{def}{=} \bigcup \{\alpha \in (A_{out} \cup \mathbb{R}_{\geq 0}) \mid s \in S. s \xrightarrow{\alpha}\}$

- Timed Input/Output Conformance relation:

$$m \text{ tioco } s \stackrel{def}{=} \forall \sigma \in TTr(s). Out(m \text{ After } \sigma) \subseteq Out(s \text{ After } \sigma)$$



- Testing in *relation* to specific assumptions about *environment*.
- Relativized Timed Input/Output Conformance:

$$m \text{ rtioco}_e s \stackrel{\text{def}}{=} \forall \sigma \in \text{TTr}(e). \text{Out}((e, m) \text{ After } \sigma) \subseteq \text{Out}((e, s) \text{ After } \sigma)$$

- Under certain conditions, rtioco is the same as *trace inclusion*:

$$m \text{ rtioco}_e s \iff \text{TTr}(m) \cap \text{TTr}(e) \subseteq \text{TTr}(s) \cap \text{TTr}(e)$$

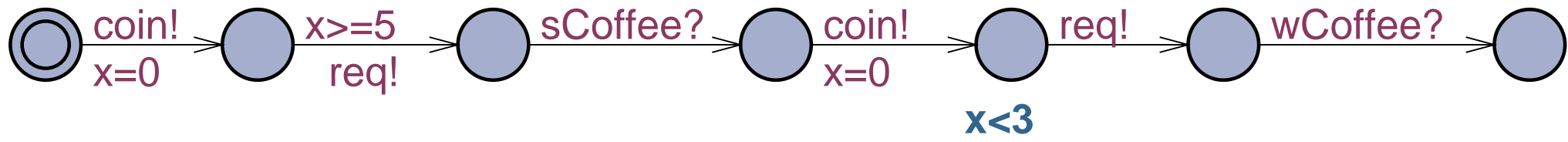
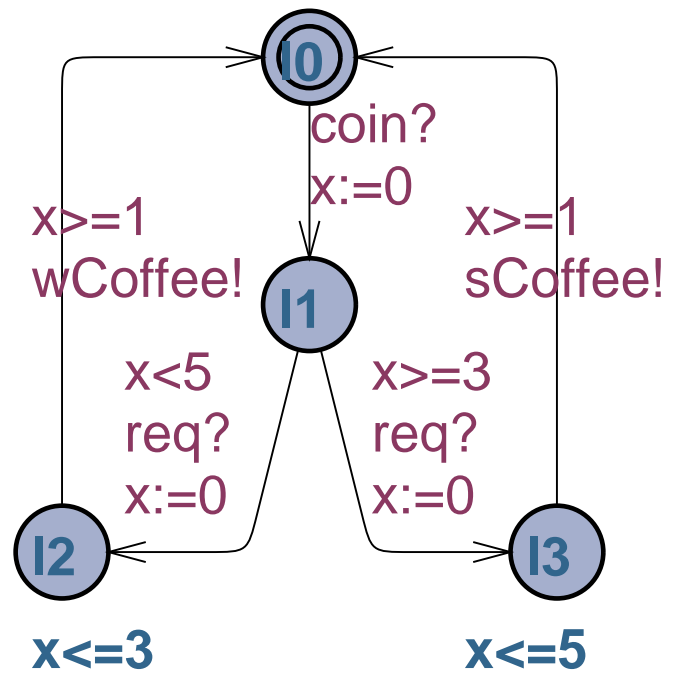
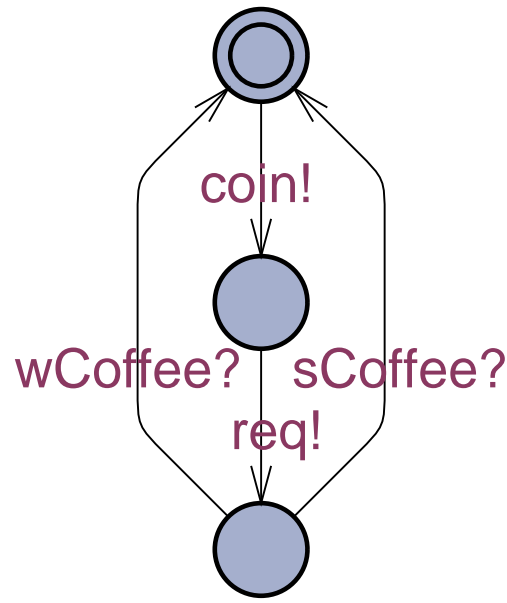
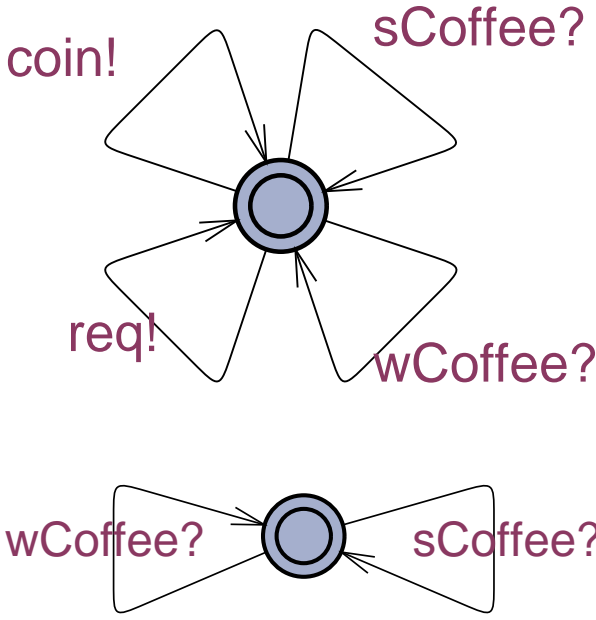
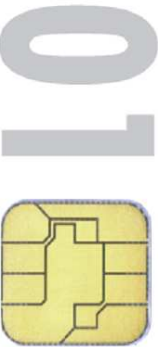
- *Ordering*: g is stronger than f :

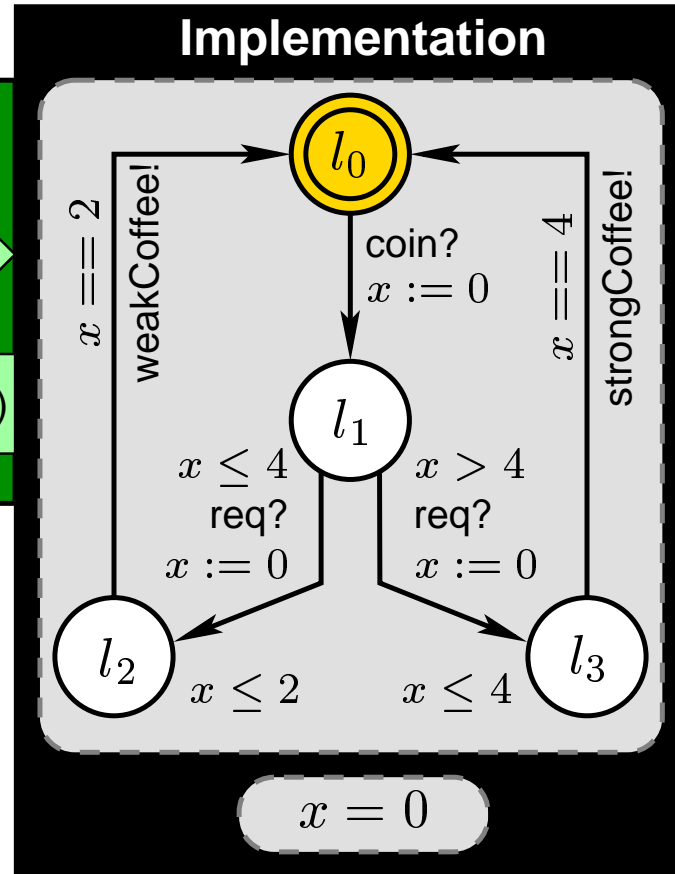
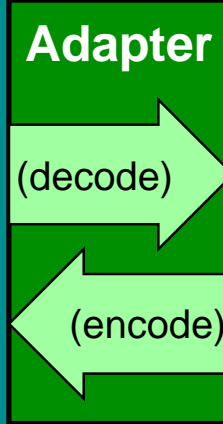
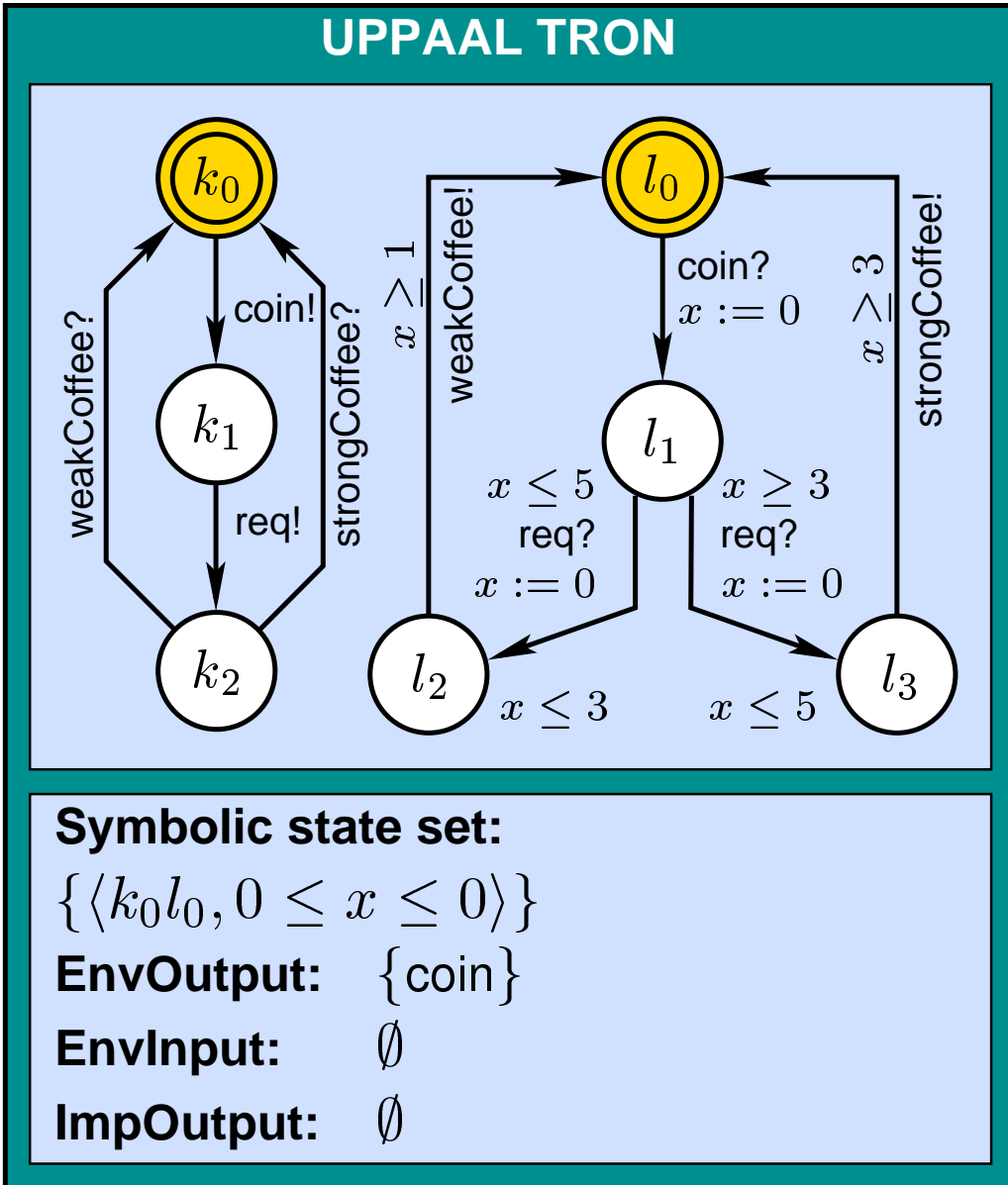
$$g \sqsubseteq f \stackrel{\text{def}}{=} \text{TTr}(g) \subseteq \text{TTr}(f)$$

- Allows to *reuse testing effort*, assuming $m \text{ rtioco}_e s$:
 - Stronger environment e' : if $e' \sqsubseteq e$ then $m \text{ rtioco}_{e'} s$.
 - Weaker specification s' : if $s \sqsubseteq s'$ then $m \text{ rtioco}_e s'$.

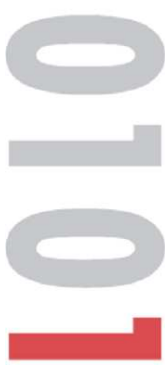
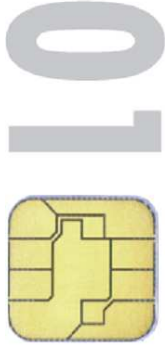


Relativized Conformance Example

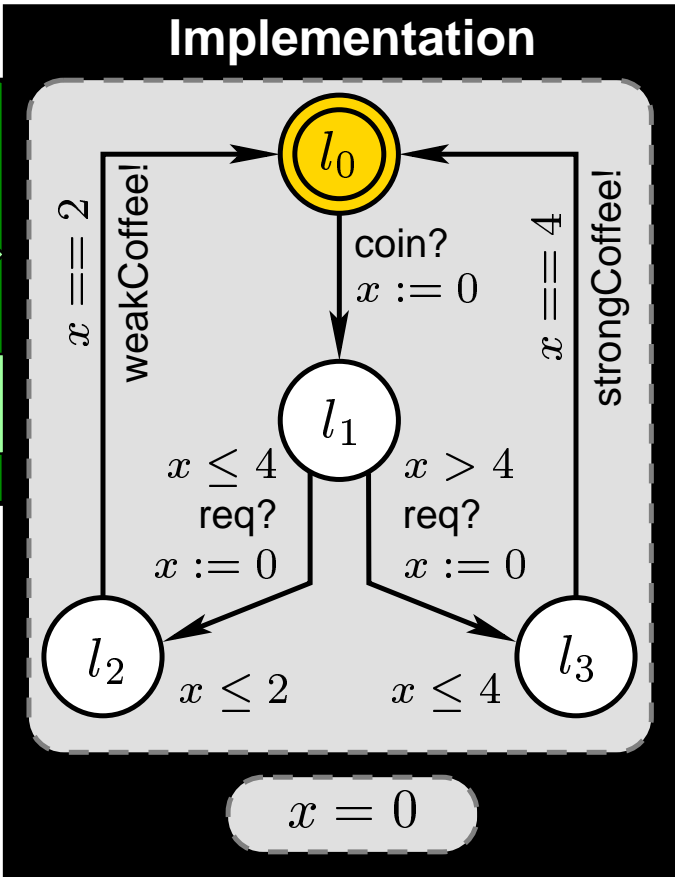
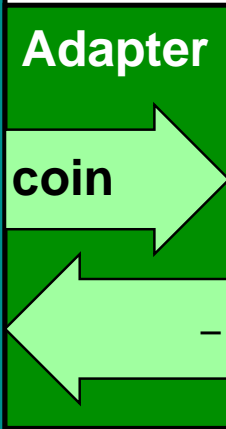
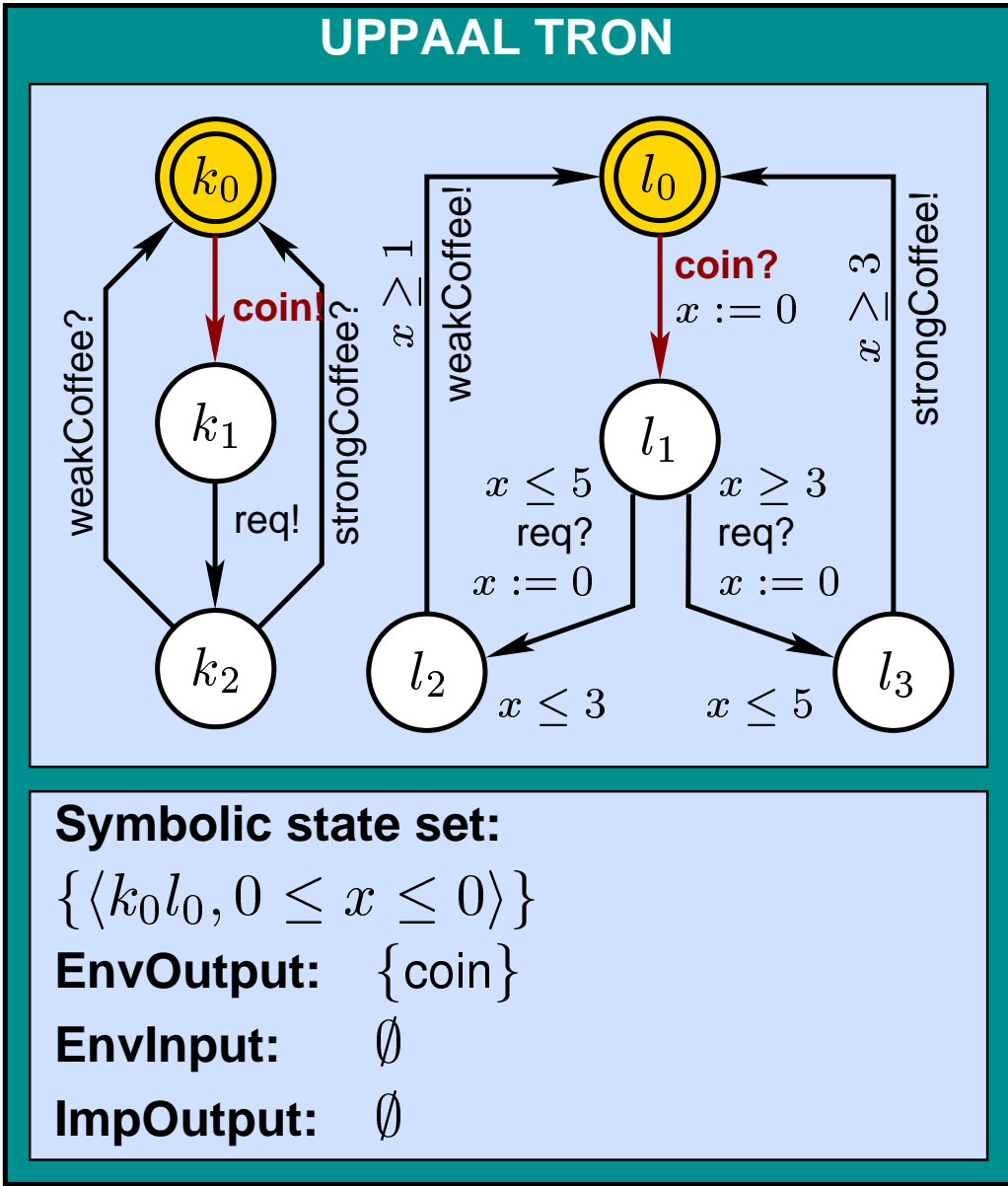




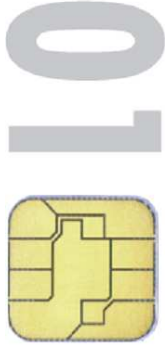
Wait for output (delay) or offer input?

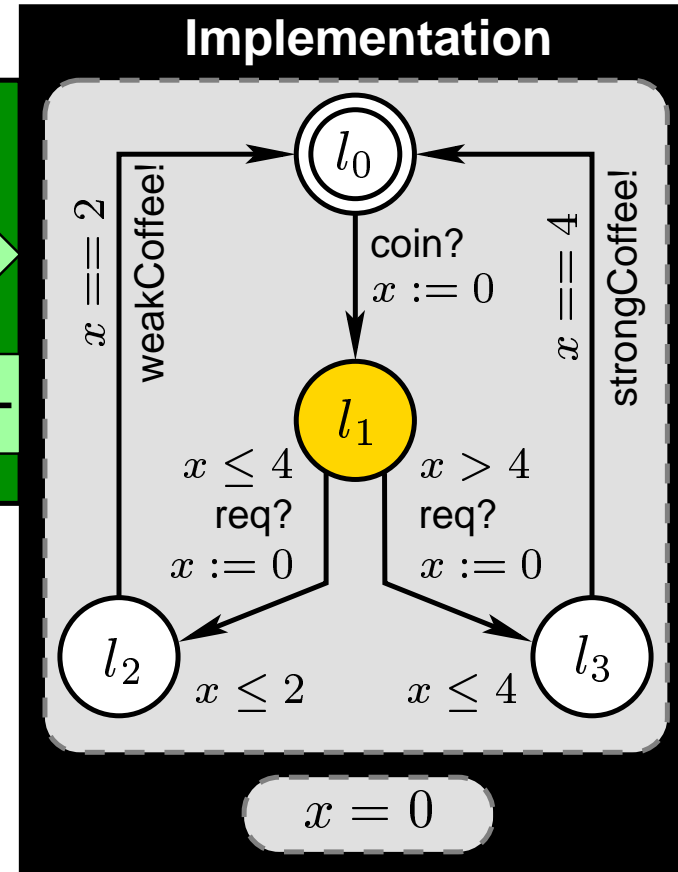
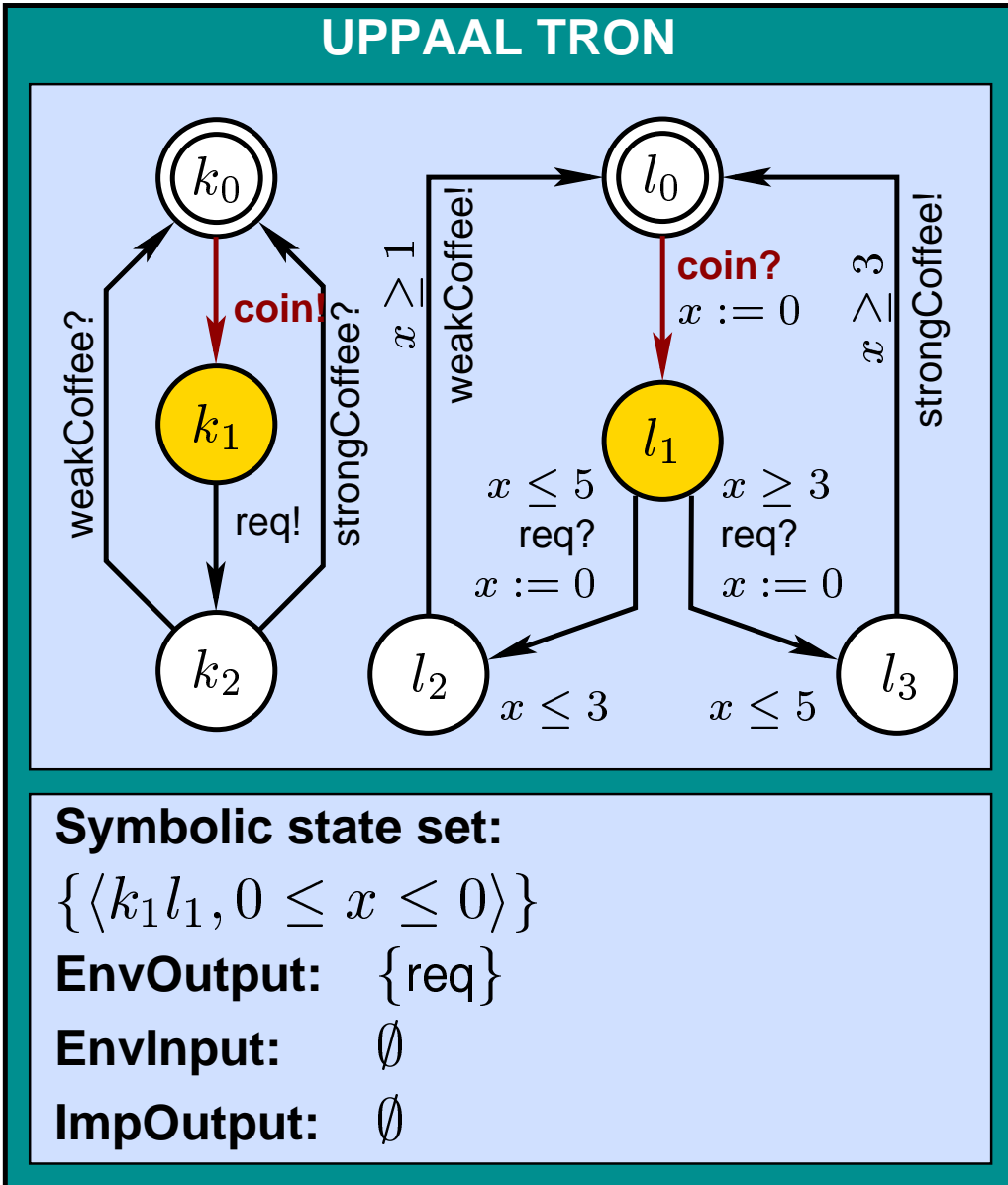


Online Test in Action

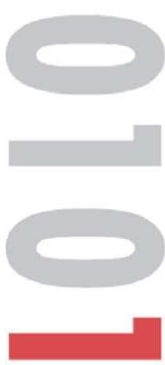
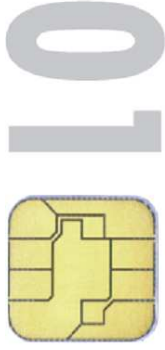


Let's offer input
choose (the only) "coin"

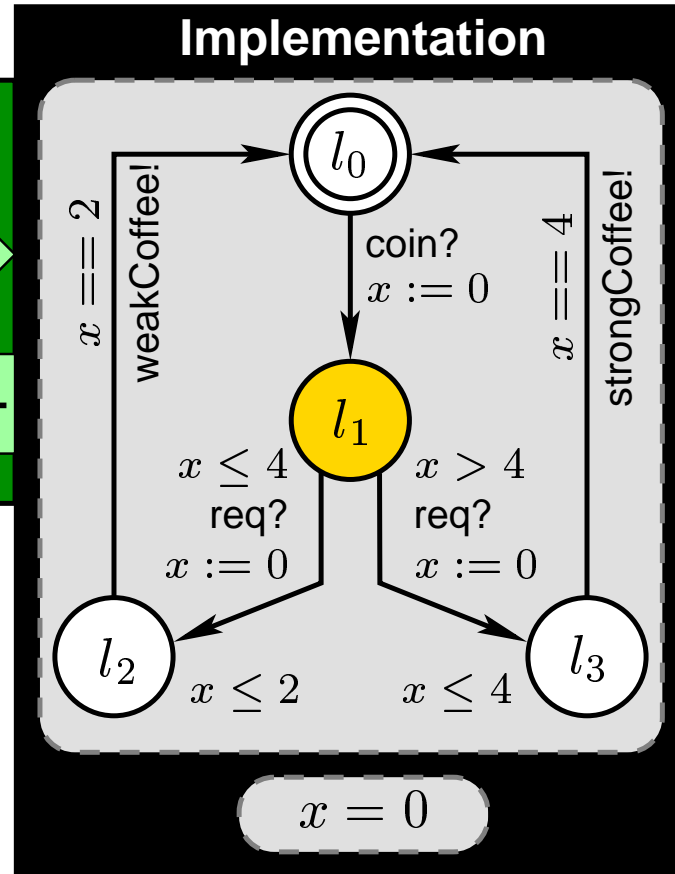
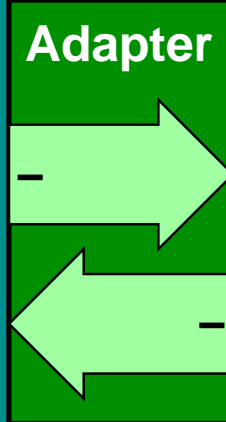
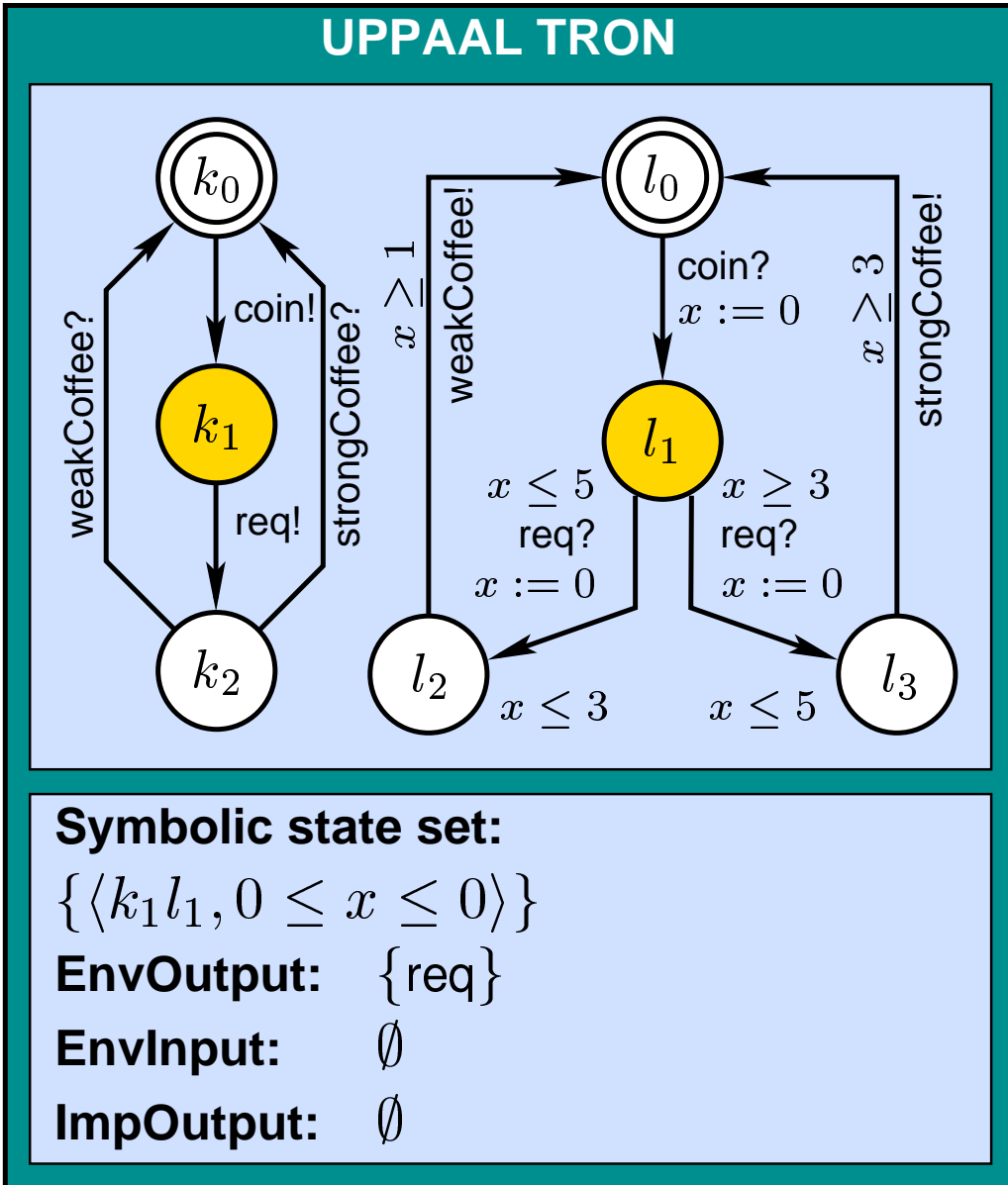




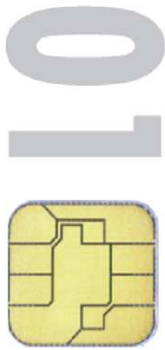
Update the state set and other variables

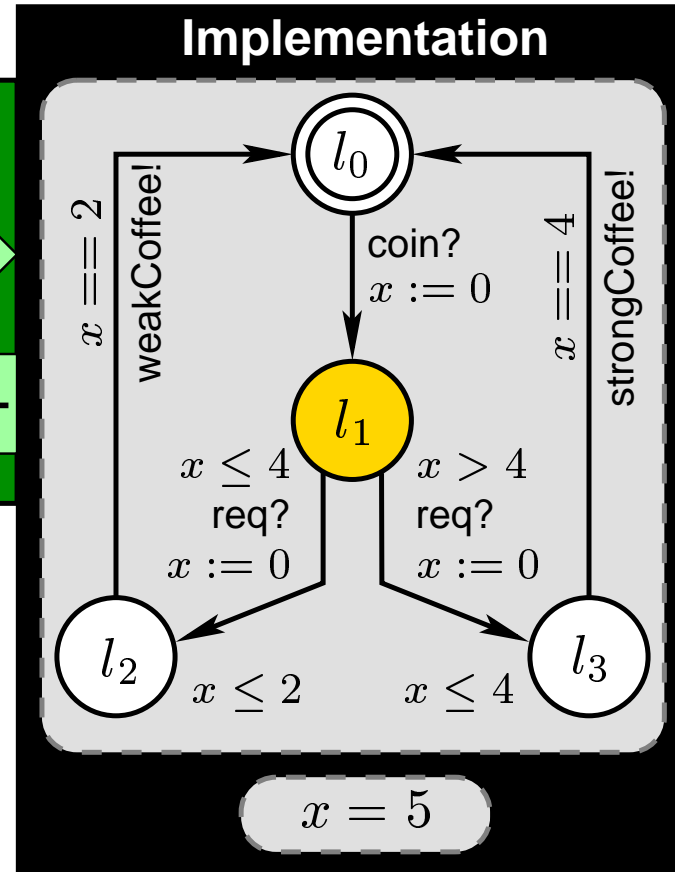
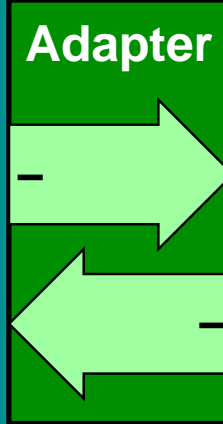
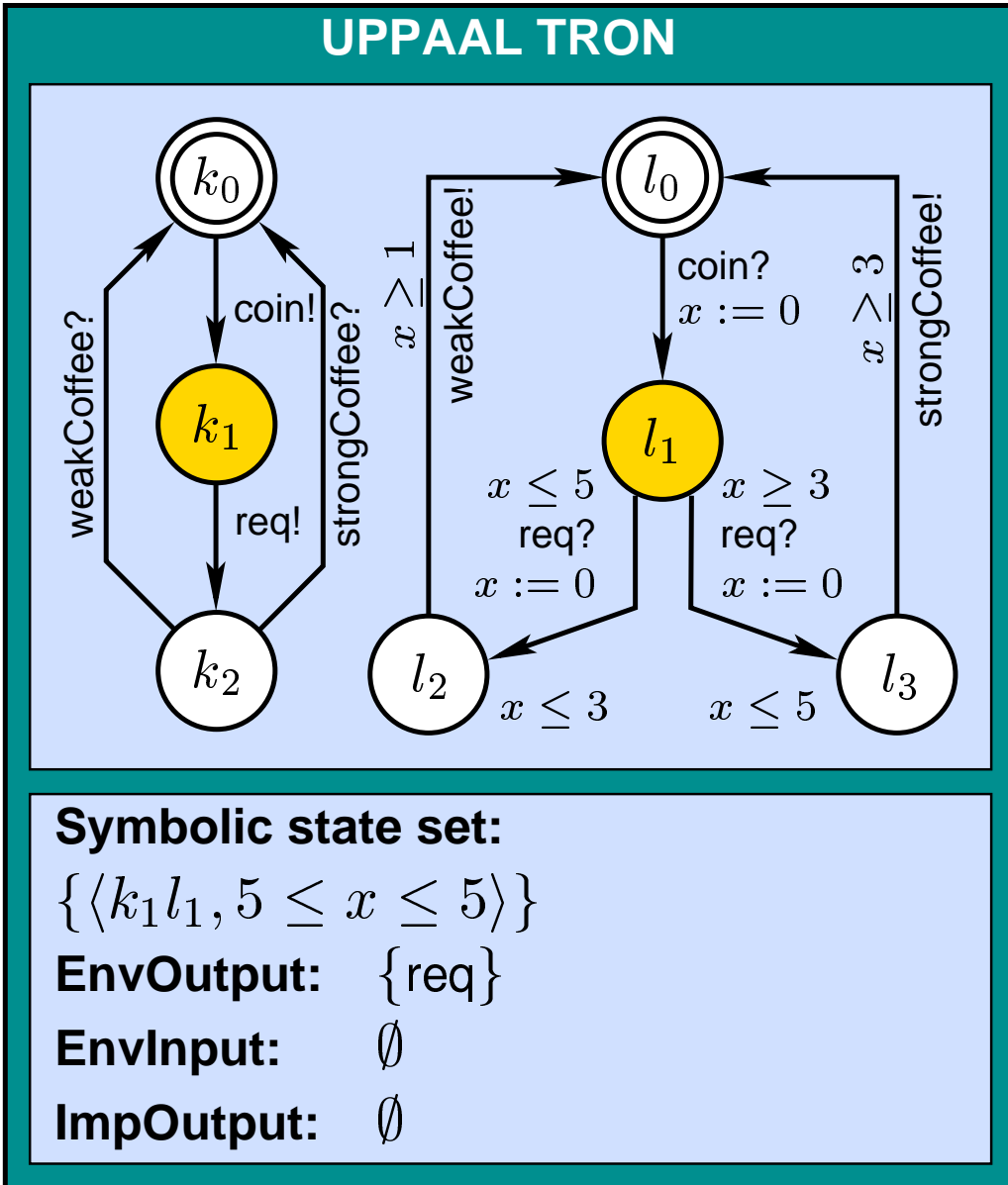


Online Test in Action

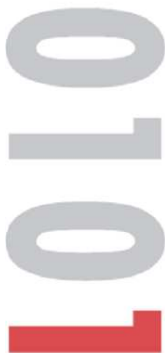
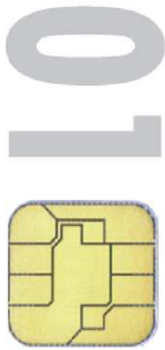


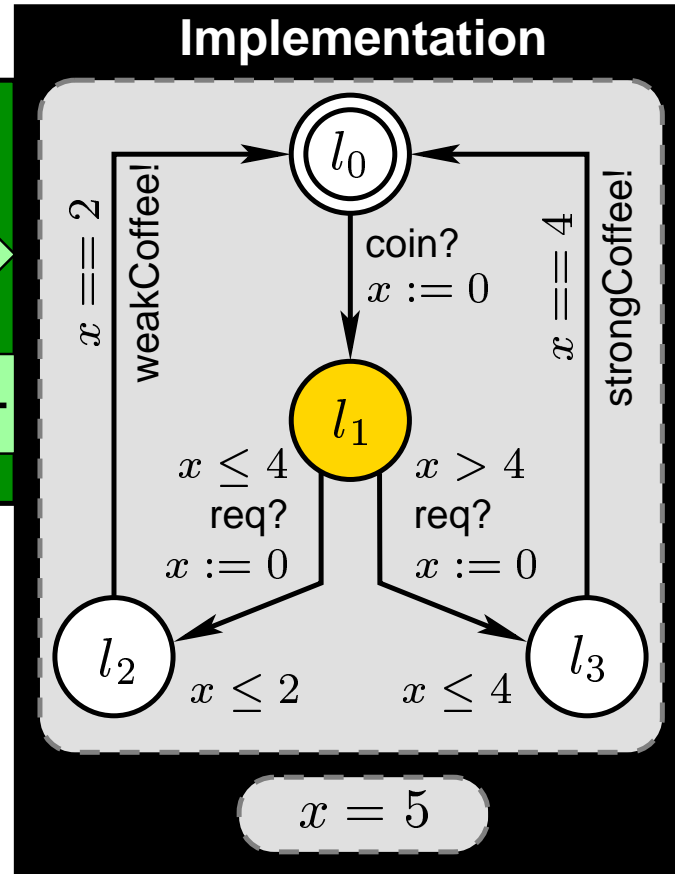
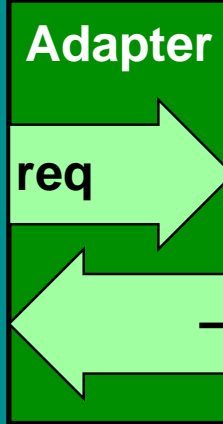
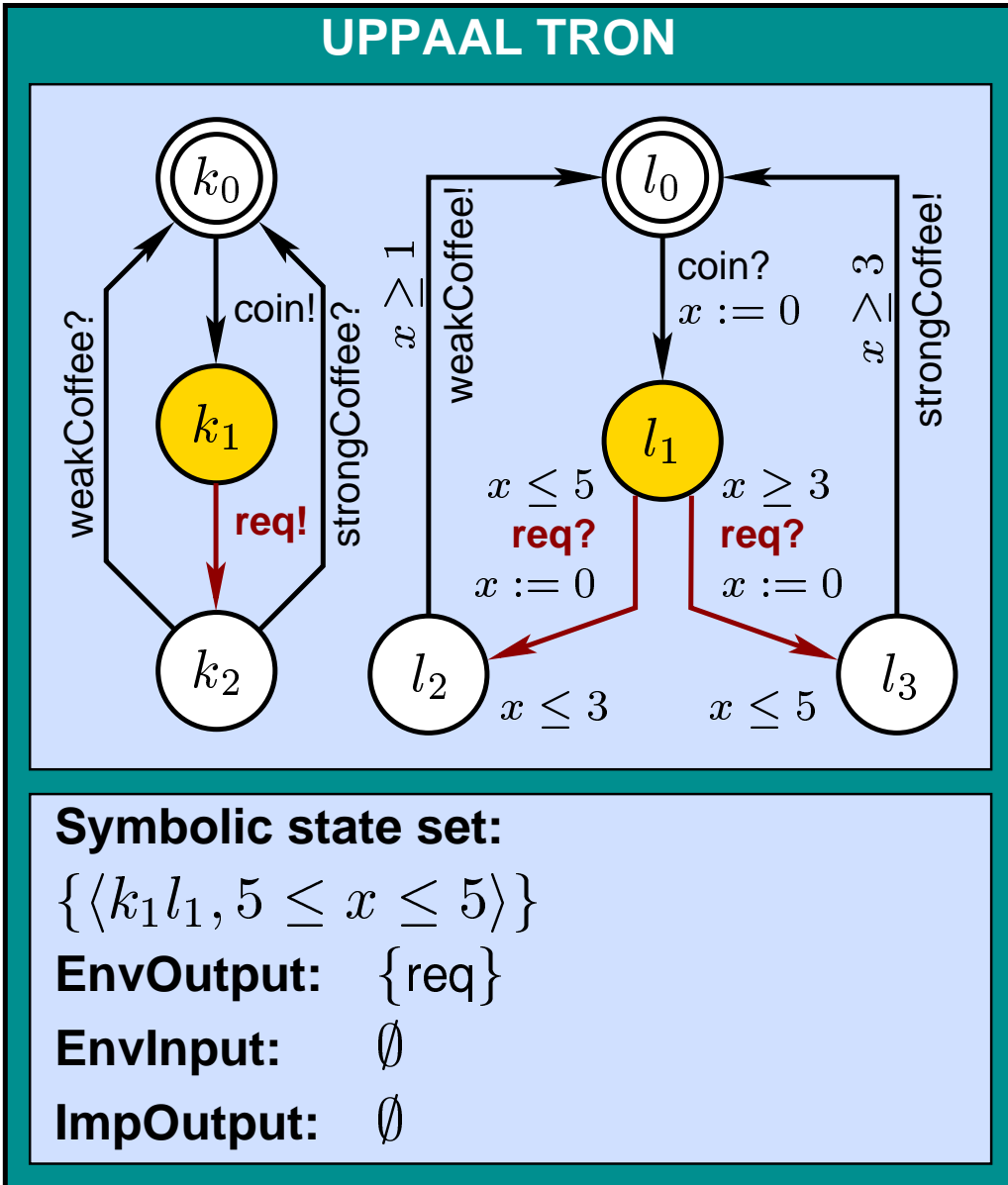
**Wait or offer input?
Let's wait for 5 units**



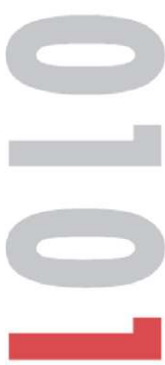
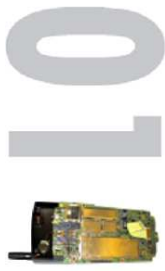
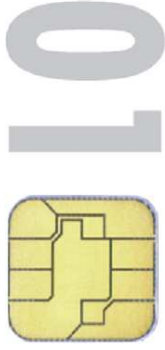


..no output so far:
update the state set..

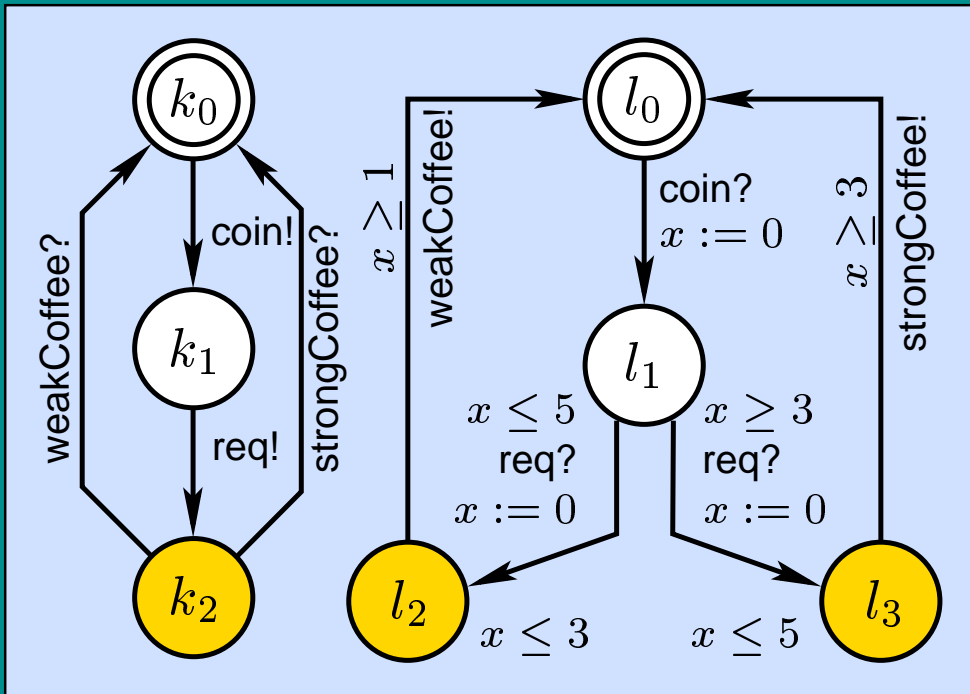




Wait or offer input?
let's offer "req"



UPPAAL TRON



Symbolic state set:

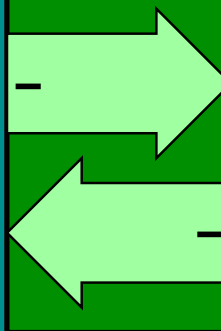
$\{\langle k_2 l_2, 0 \leq x \leq 0 \rangle, \langle k_2 l_3, 0 \leq x \leq 0 \rangle\}$

EnvOutput: \emptyset

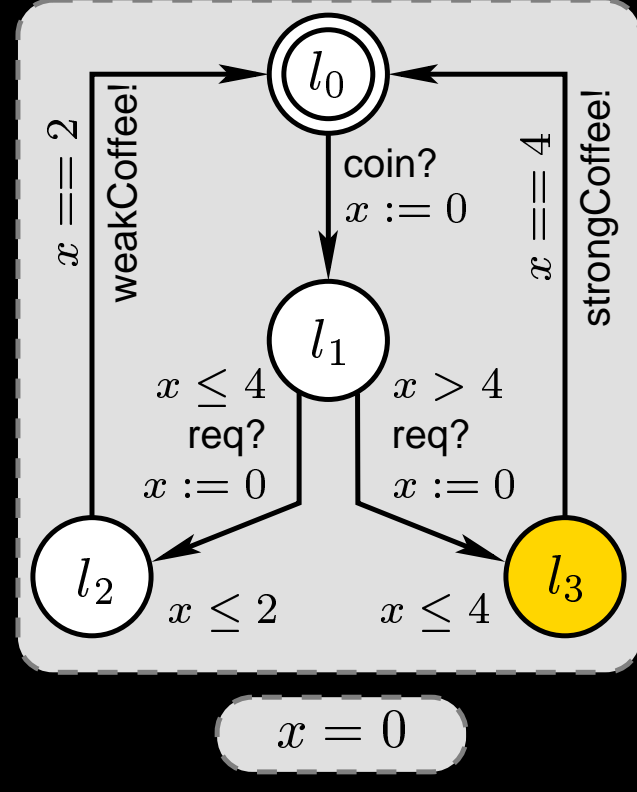
EnvInput: $\{\text{weakCoffee}, \text{strongCoffee}\}$

ImpOutput: $\{\text{weakCoffee}, \text{strongCoffee}\}$

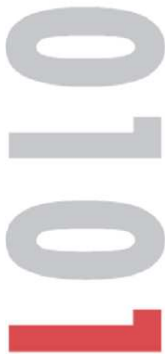
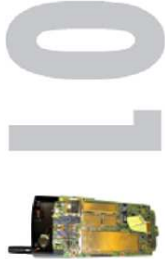
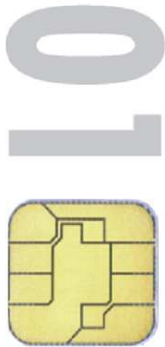
Adapter

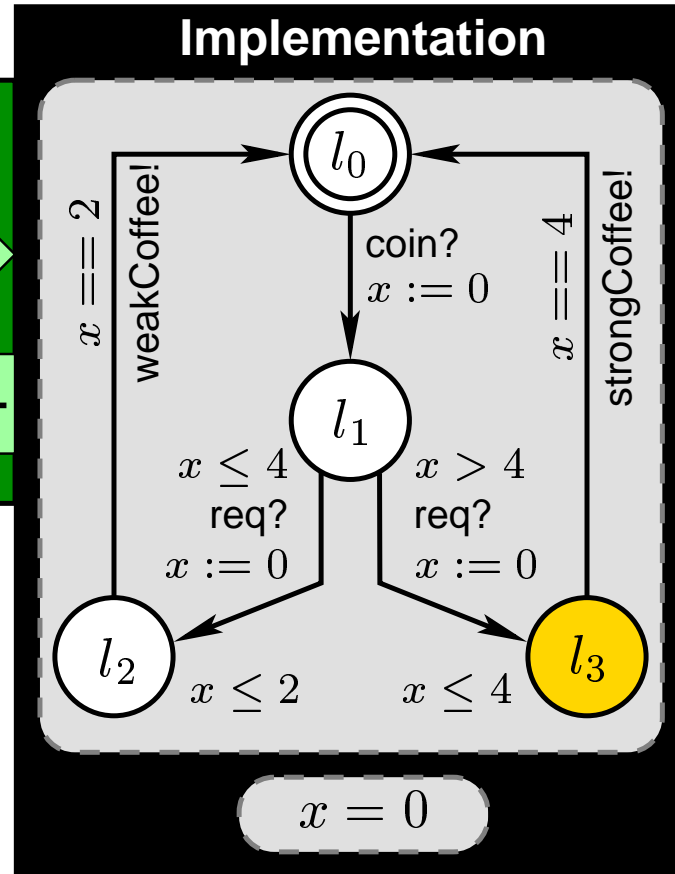
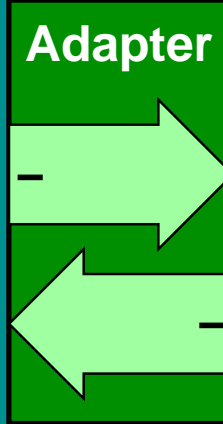
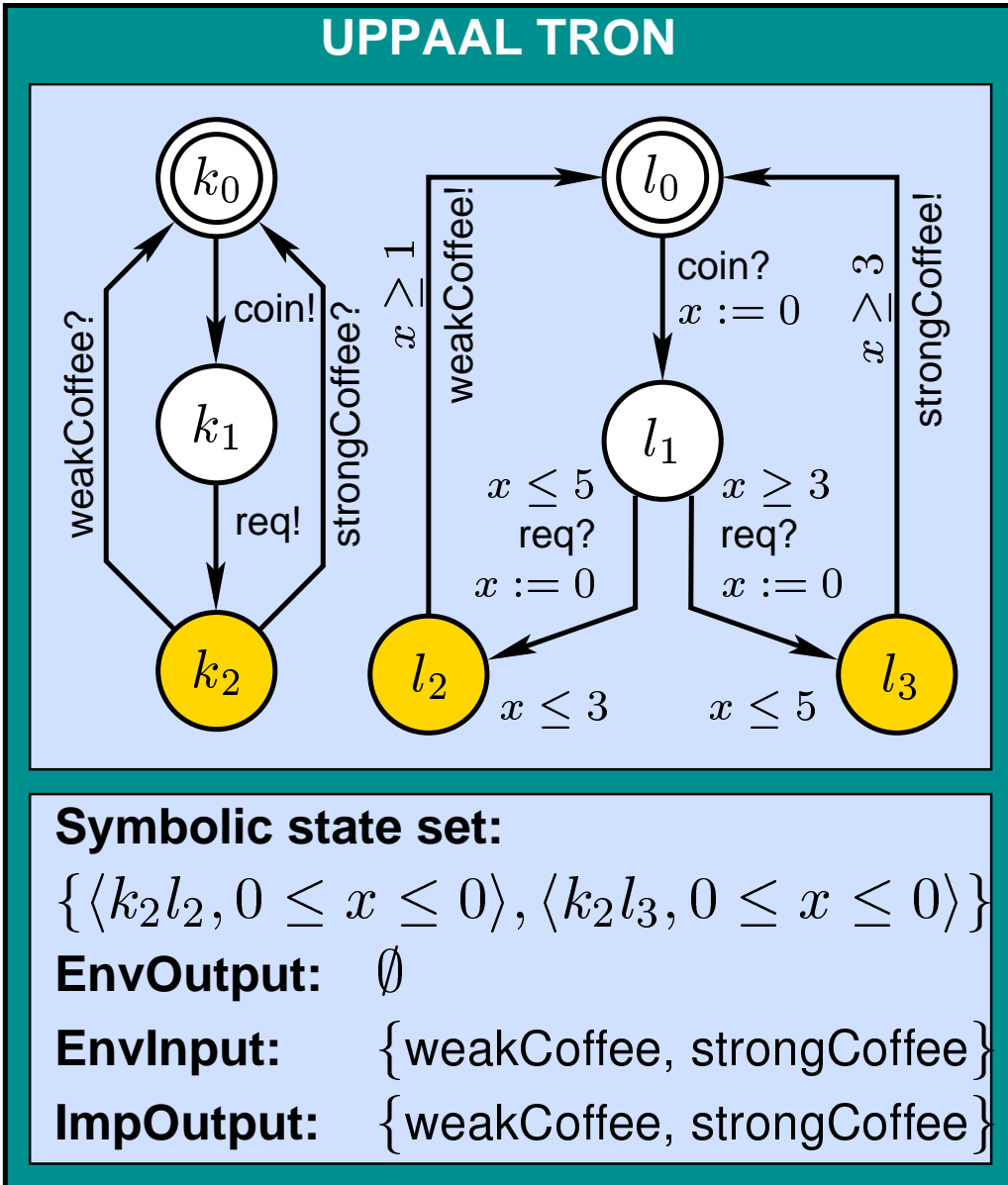


Implementation

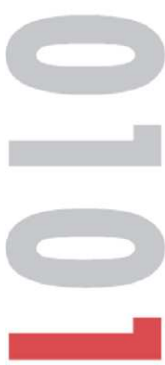
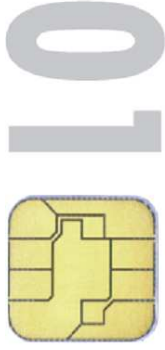


Update the state set and other variables

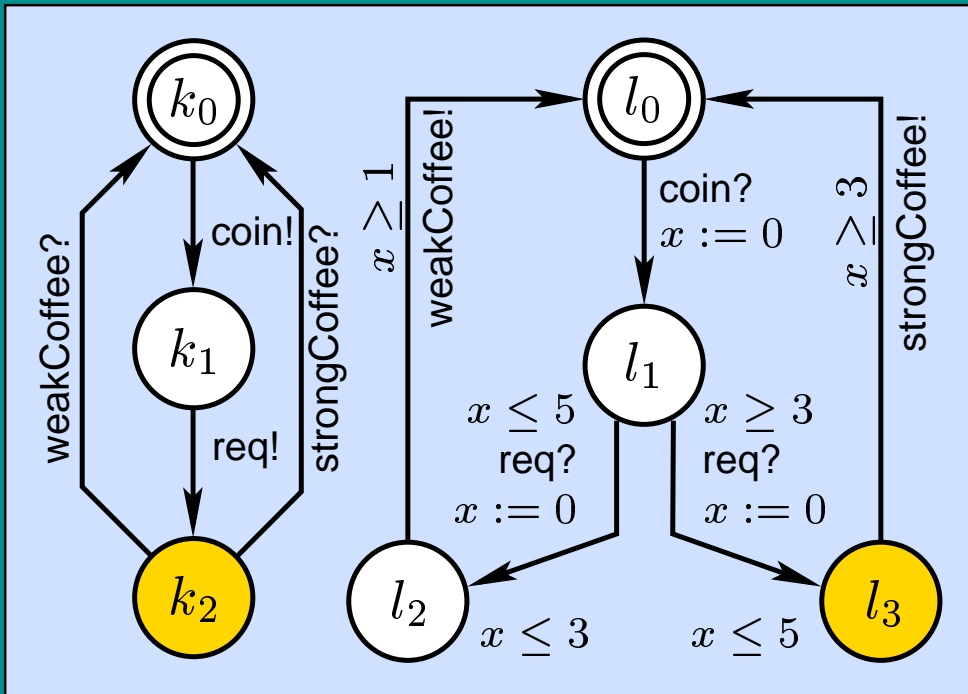




**Wait or offer input?
Let's wait for 4 units**



UPPAAL TRON



Symbolic state set:

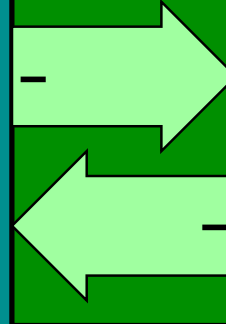
$\{\langle k_2 l_3, 4 \leq x \leq 4 \rangle\}$

EnvOutput: \emptyset

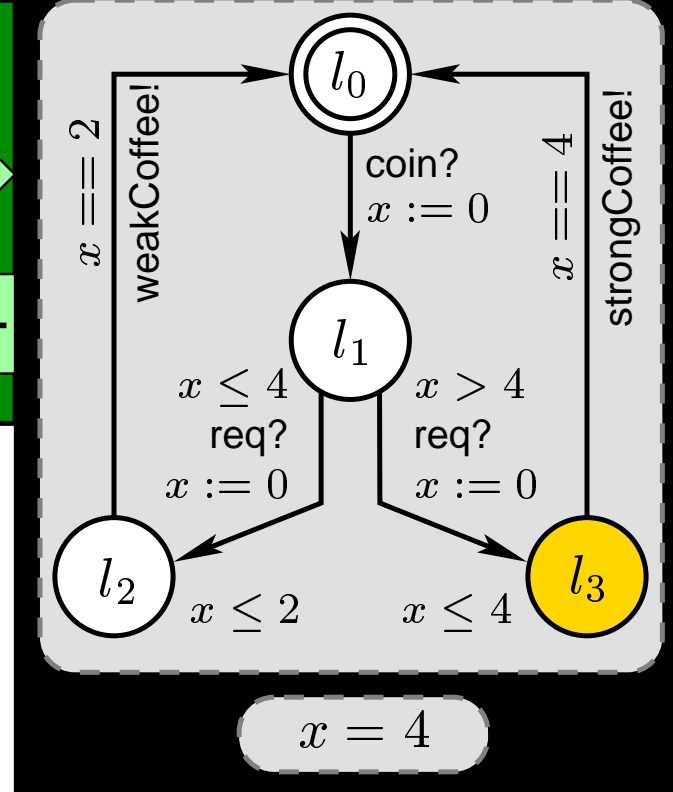
EnvInput: $\{\text{strongCoffee}\}$

ImpOutput: $\{\text{strongCoffee}\}$

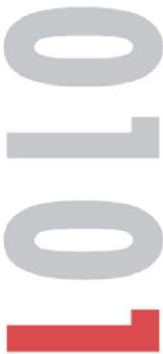
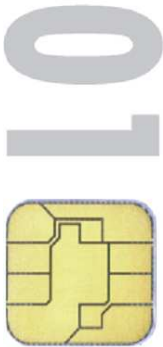
Adapter



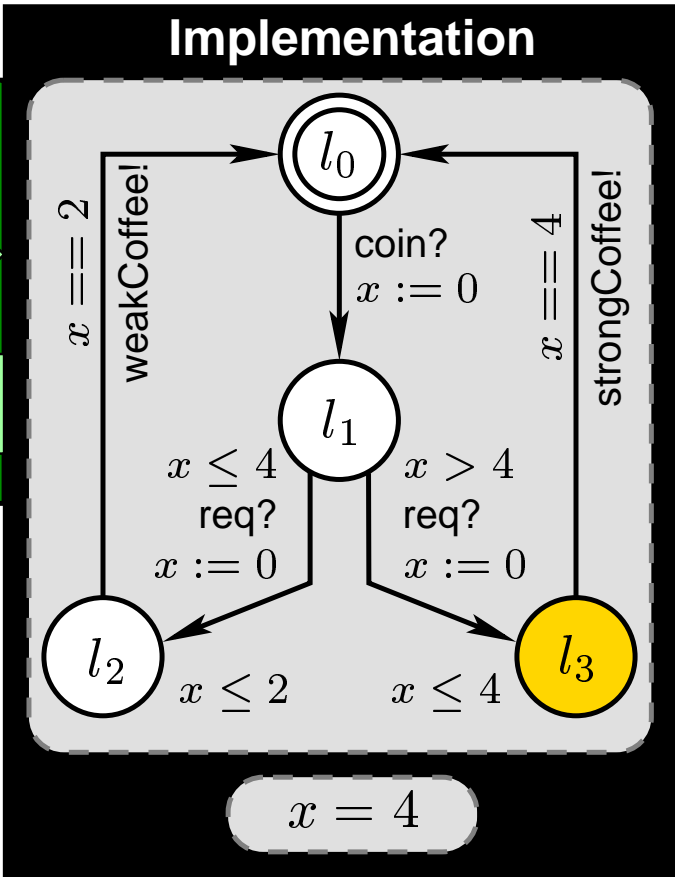
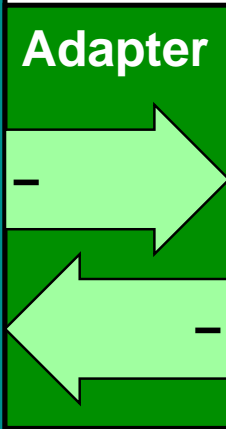
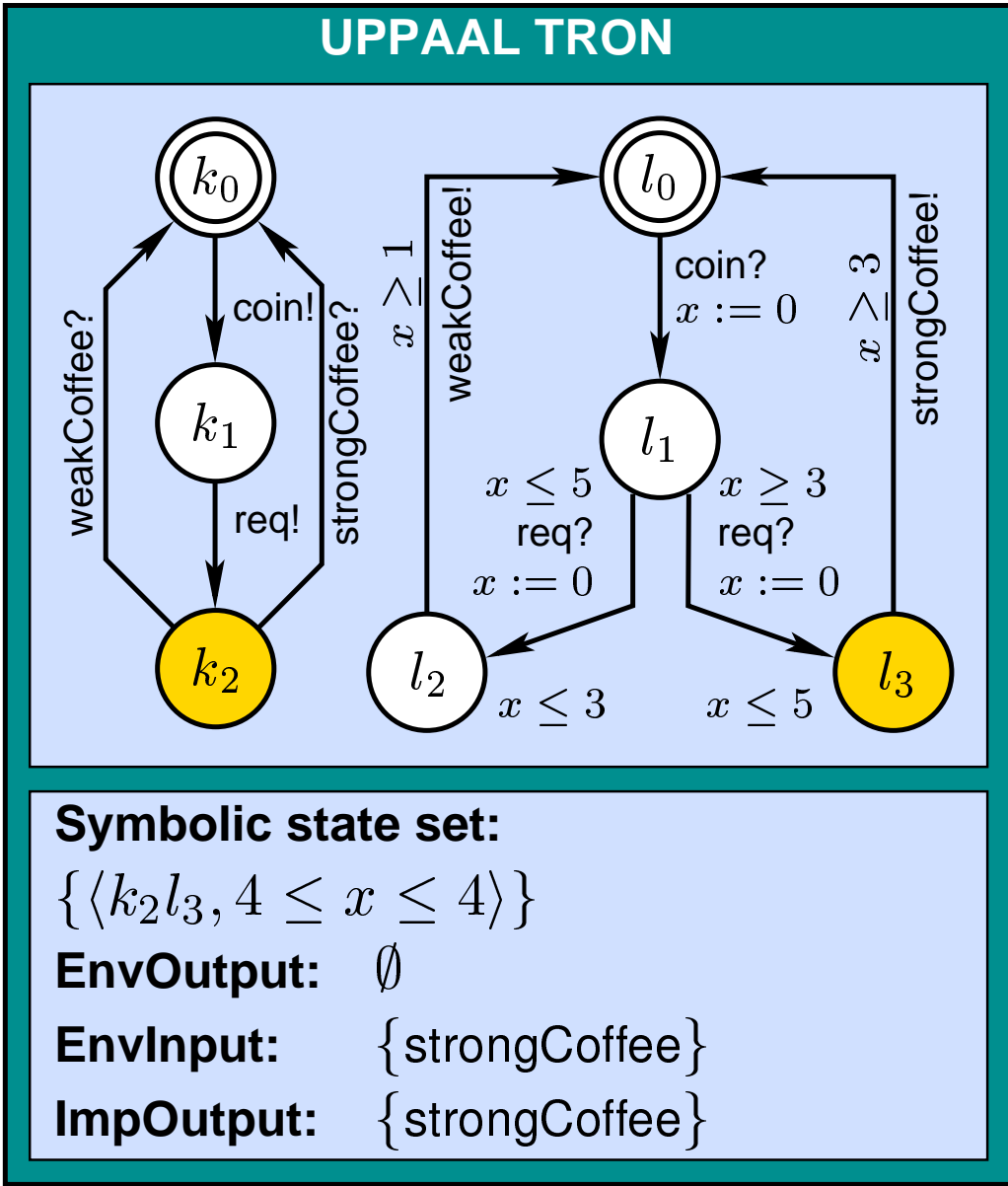
Implementation



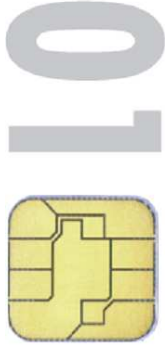
..no output so far:
update the state set..

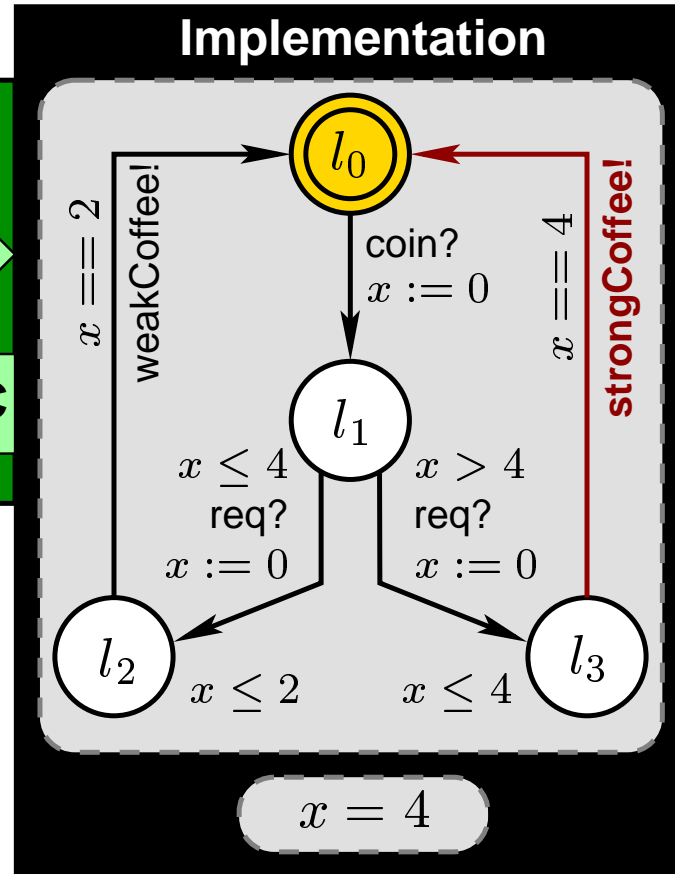
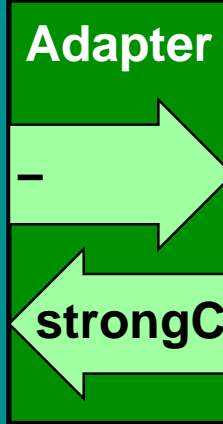
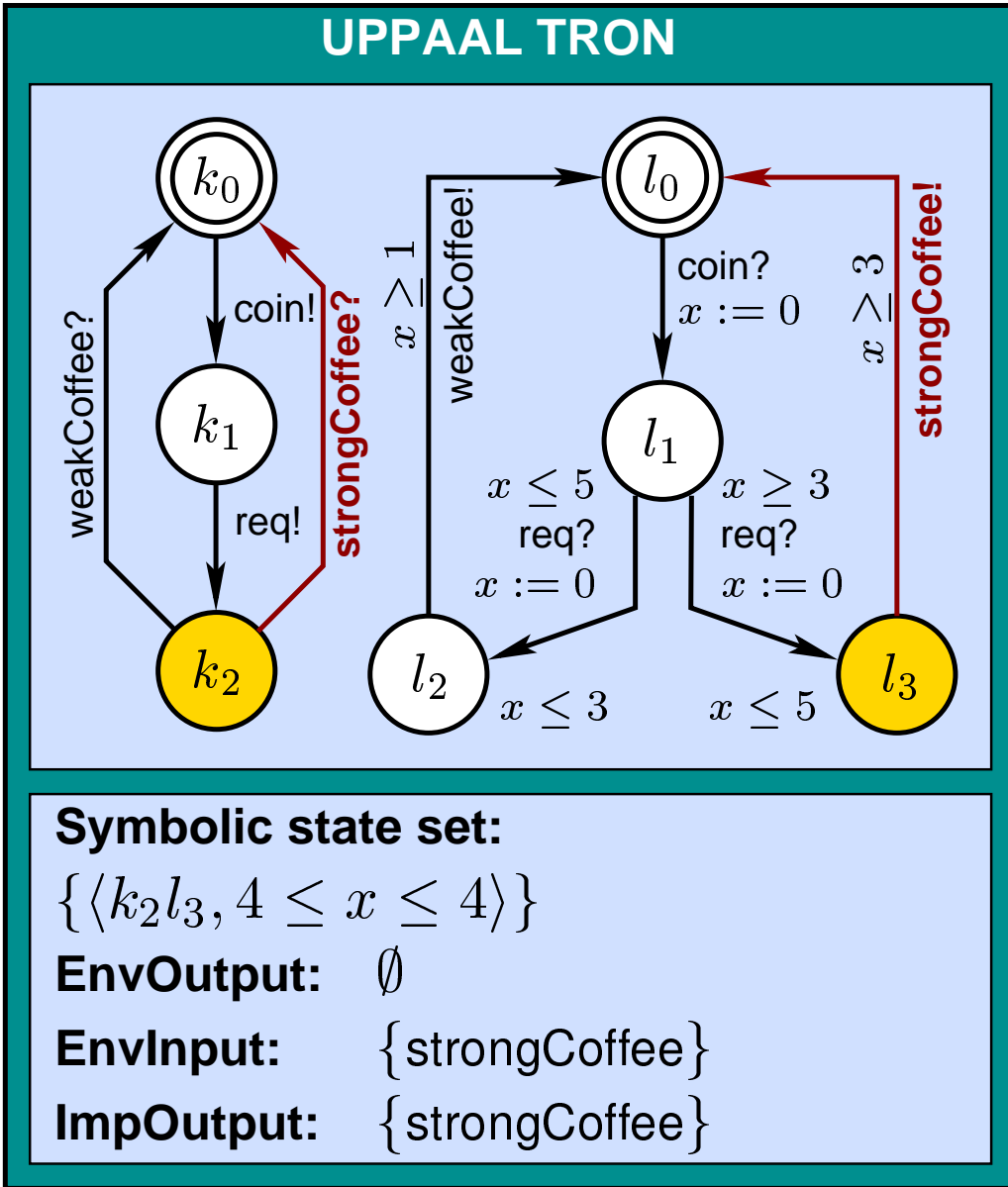


Online Test in Action

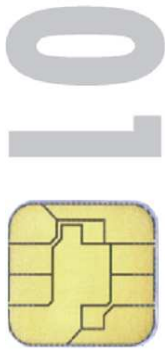


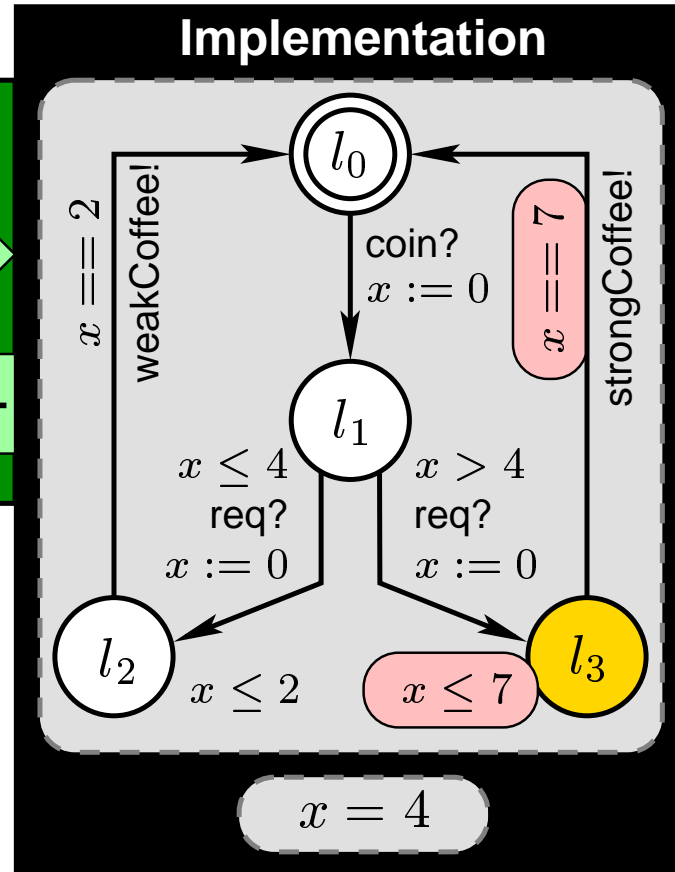
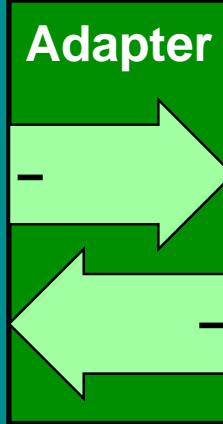
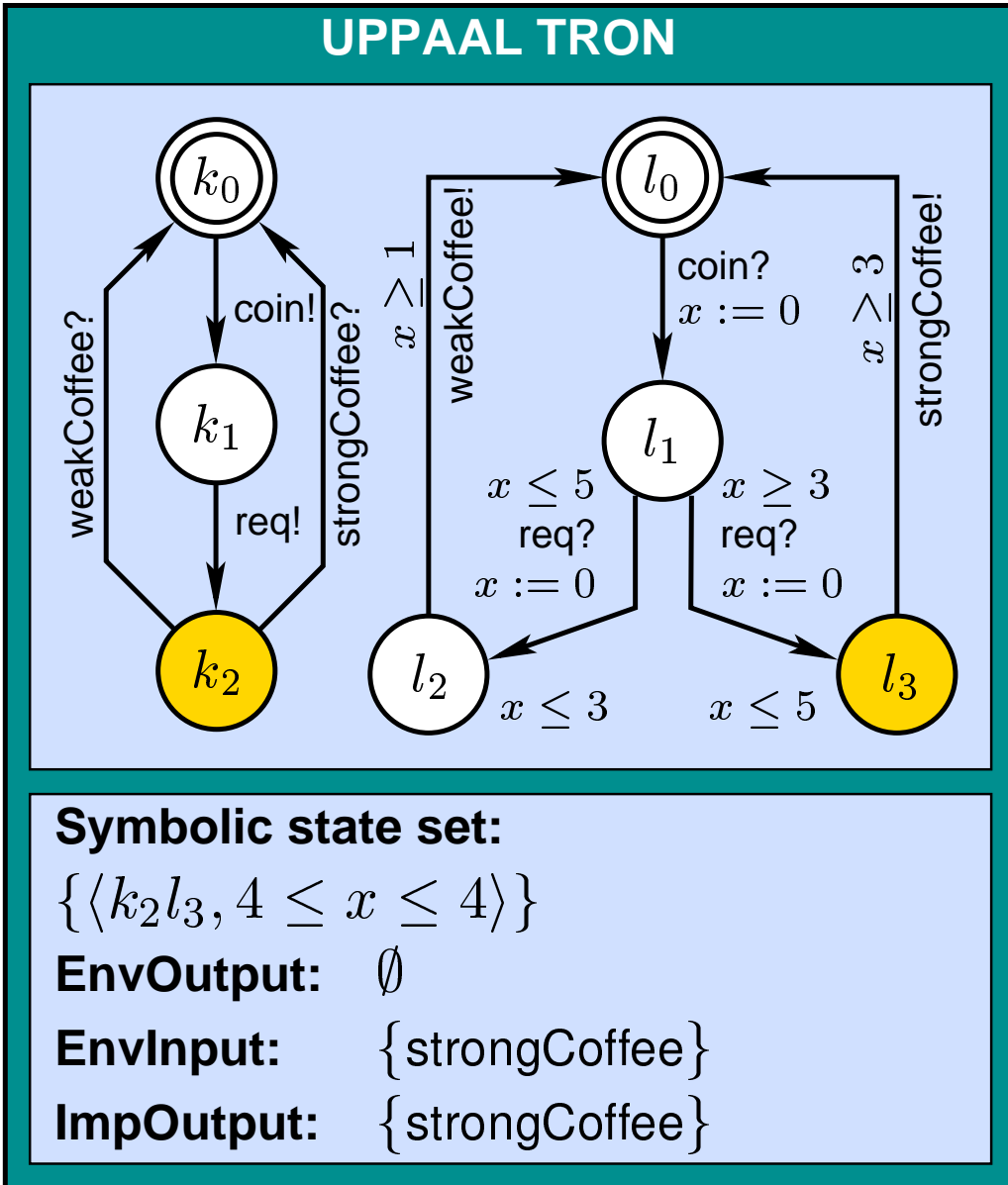
**Wait or offer input?
Let's wait for 2 units**



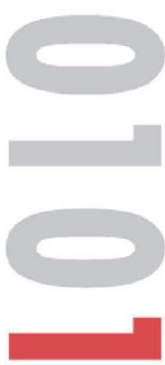
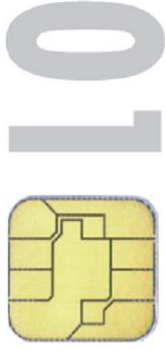


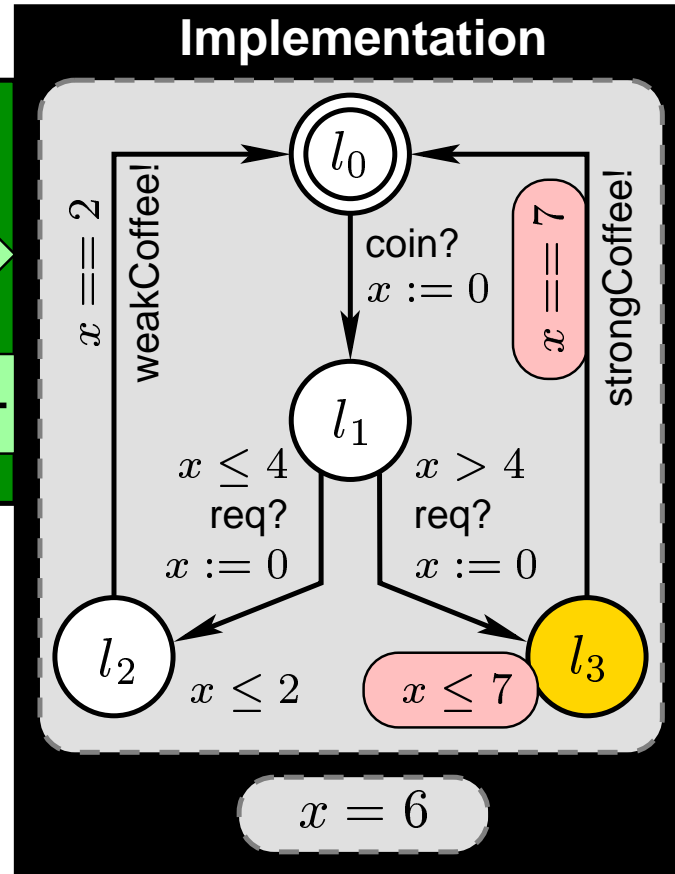
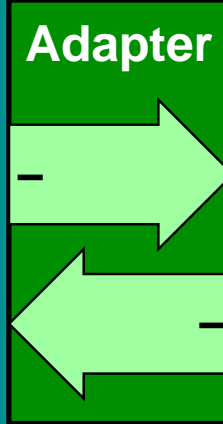
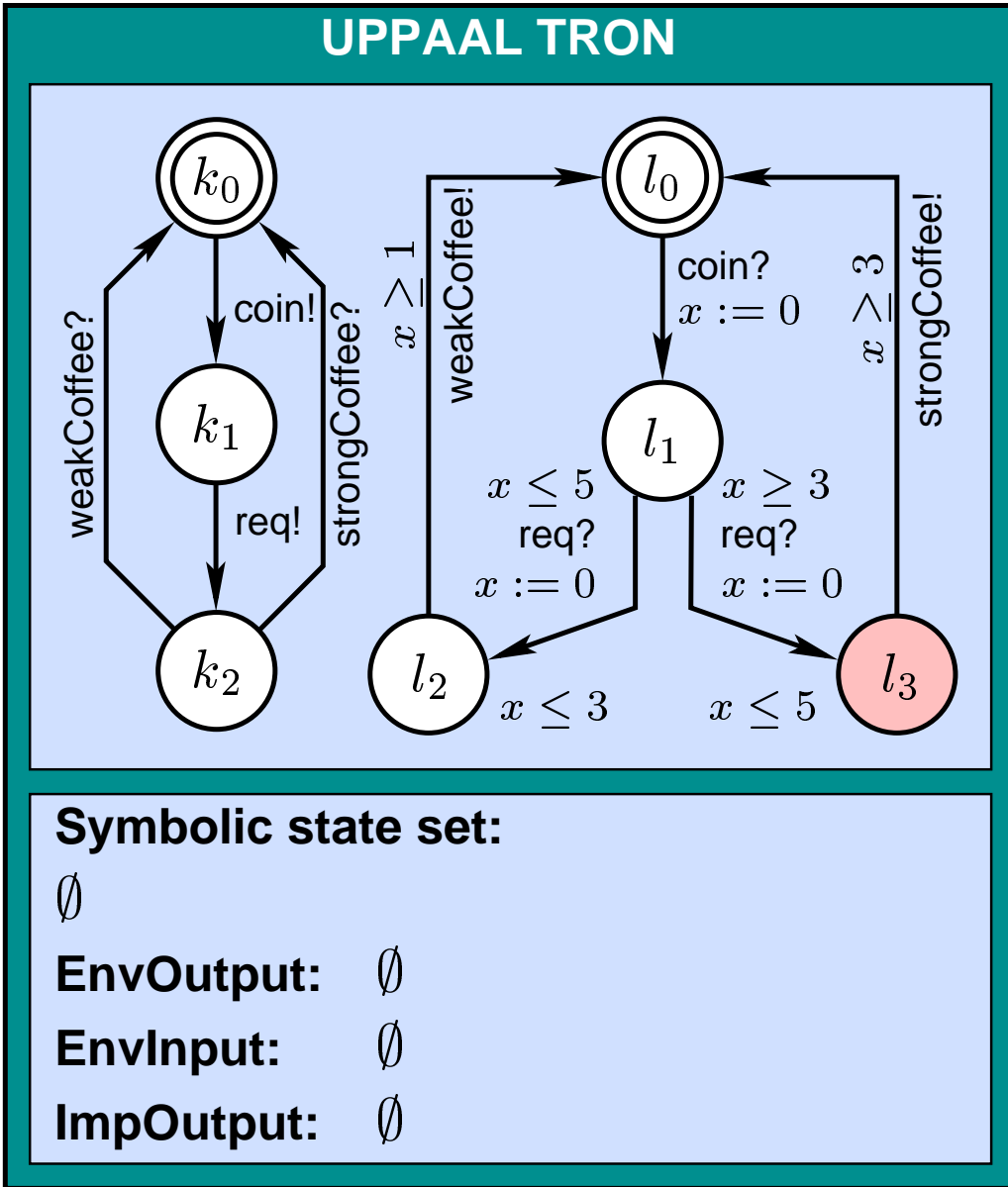
got output after 0 delay:
update the state set



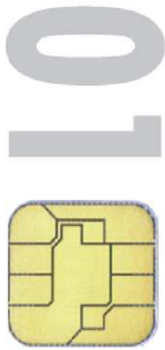


(what if there is a bug?)
 Let's wait for 2 units





..no output so far:
update the state set.. (!)



Online Test Algorithm

while $\mathcal{Z} \neq \emptyset \wedge \#iterations \leq T$ **do** choose randomly:

1. **if** $EnvOutput(\mathcal{Z}) \neq \emptyset$ // offer an input

randomly choose $a \in EnvOutput(\mathcal{Z})$

send a to IUT

$\mathcal{Z} := \mathcal{Z}$ After a

2. randomly choose $\delta \in Delays(\mathcal{Z})$ // wait for an output

sleep for δ time units and wake up on output o

if o occurs at $\delta' \leq \delta$ **then**

$\mathcal{Z} := \mathcal{Z}$ After δ'

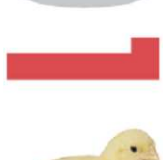
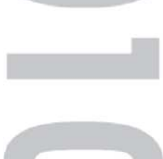
if $o \notin ImpOutput(\mathcal{Z})$ **then return fail**

else $\mathcal{Z} := \mathcal{Z}$ After o

else $\mathcal{Z} := \mathcal{Z}$ After δ // no output within δ delay

3. $\mathcal{Z} := \{(s_0, e_0)\}$, **reset** IUT //reset and restart

if $\mathcal{Z} = \emptyset$ **then return fail else return pass**



Online Test Algorithm

while $\mathcal{Z} \neq \emptyset \wedge \#iterations \leq T$ **do** choose randomly:

1. **if** $EnvOutput(\mathcal{Z}) \neq \emptyset$

// offer an input

randomly choose $a \in EnvOutput(\mathcal{Z})$

send a to IUT

$\mathcal{Z} := \mathcal{Z}$ After a

2. randomly choose $\delta \in Delays(\mathcal{Z})$

// wait for an output

sleep for δ time units and wake up on output o

if o occurs at $\delta' \leq \delta$ **then**

$\mathcal{Z} := \mathcal{Z}$ After δ'

if $o \notin ImpOutput(\mathcal{Z})$ **then return fail**

else $\mathcal{Z} := \mathcal{Z}$ After o

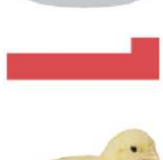
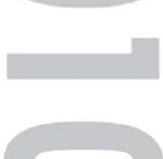
else $\mathcal{Z} := \mathcal{Z}$ After δ

// no output within δ delay

// reset and restart

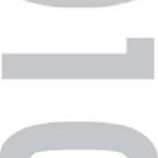
3. $\mathcal{Z} := \{(s_0, e_0)\}$, **reset** IUT

if $\mathcal{Z} = \emptyset$ **then return fail** **else return pass**

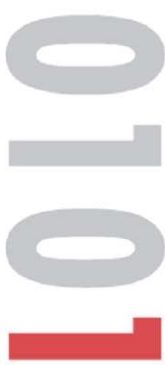
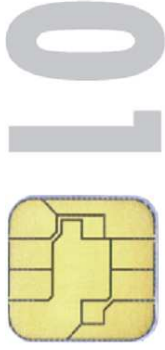
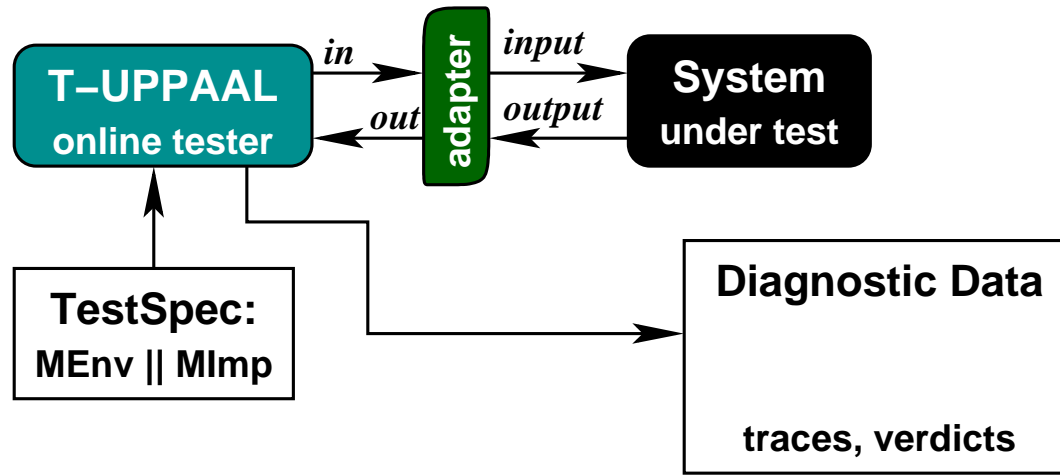


sound and complete in limit

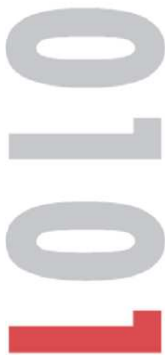
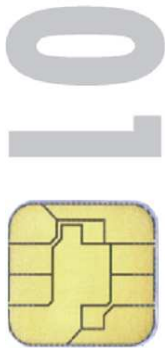
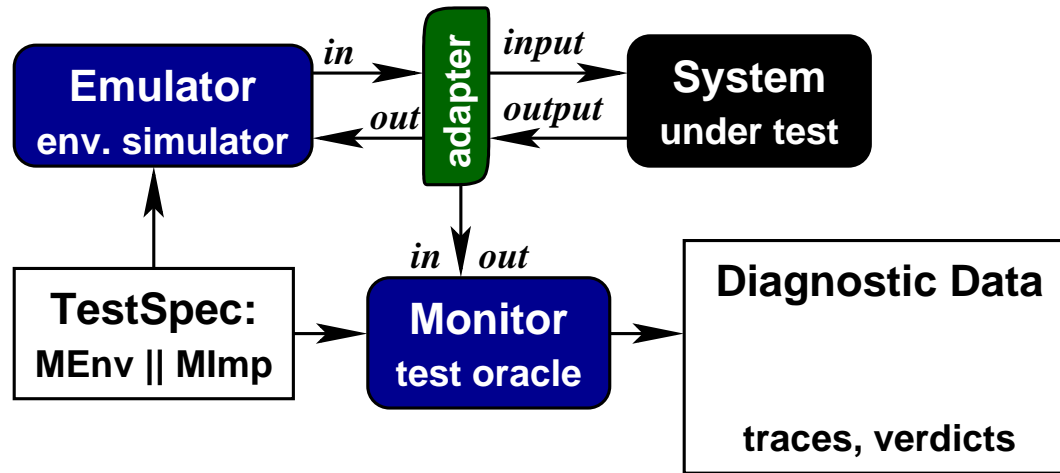
- Online real-time testing theoretically *sound* and *complete*.
- Relativized conformance provides:
 - Explicit environment assumptions
 - Realism and guiding
 - Separation of concerns
 - Modularity
 - Theoretical properties
- Implementation:
 - Allows *abstract* and *non-deterministic* specifications
 - Shows encouraging *error detection* capability and *performance*
 - Using UPPAAL model-checker constructs, online testing setup can also be applied to other model-checker.



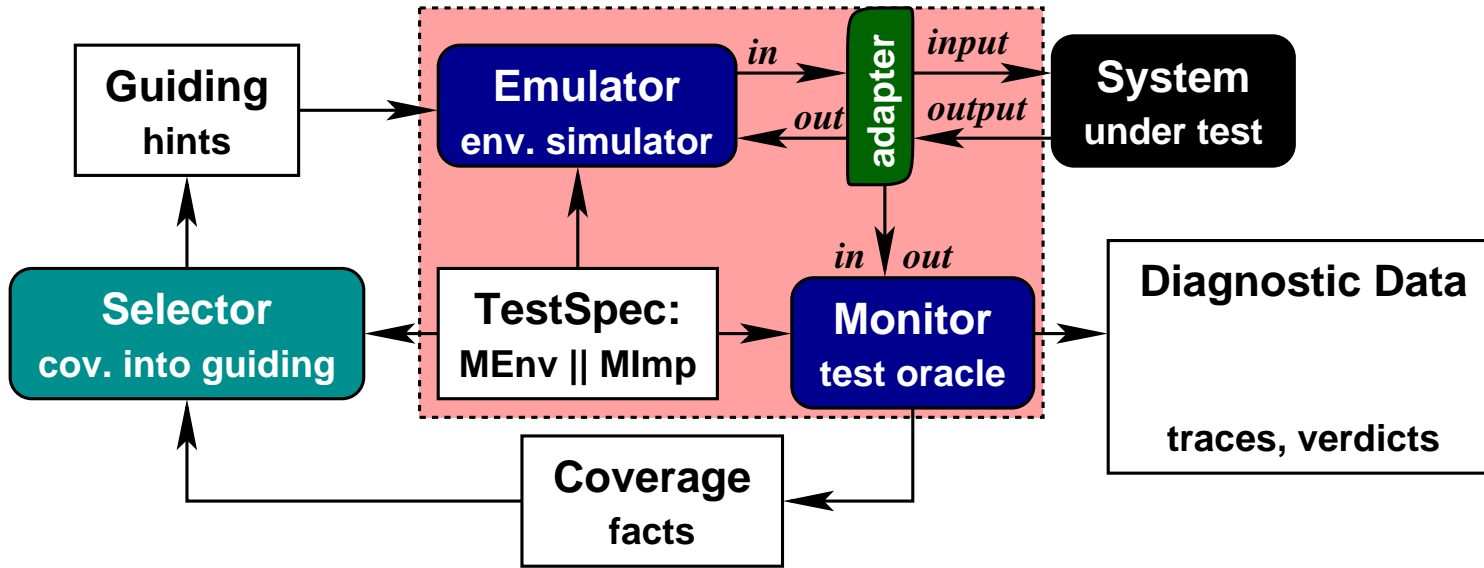
Summary and Future Work



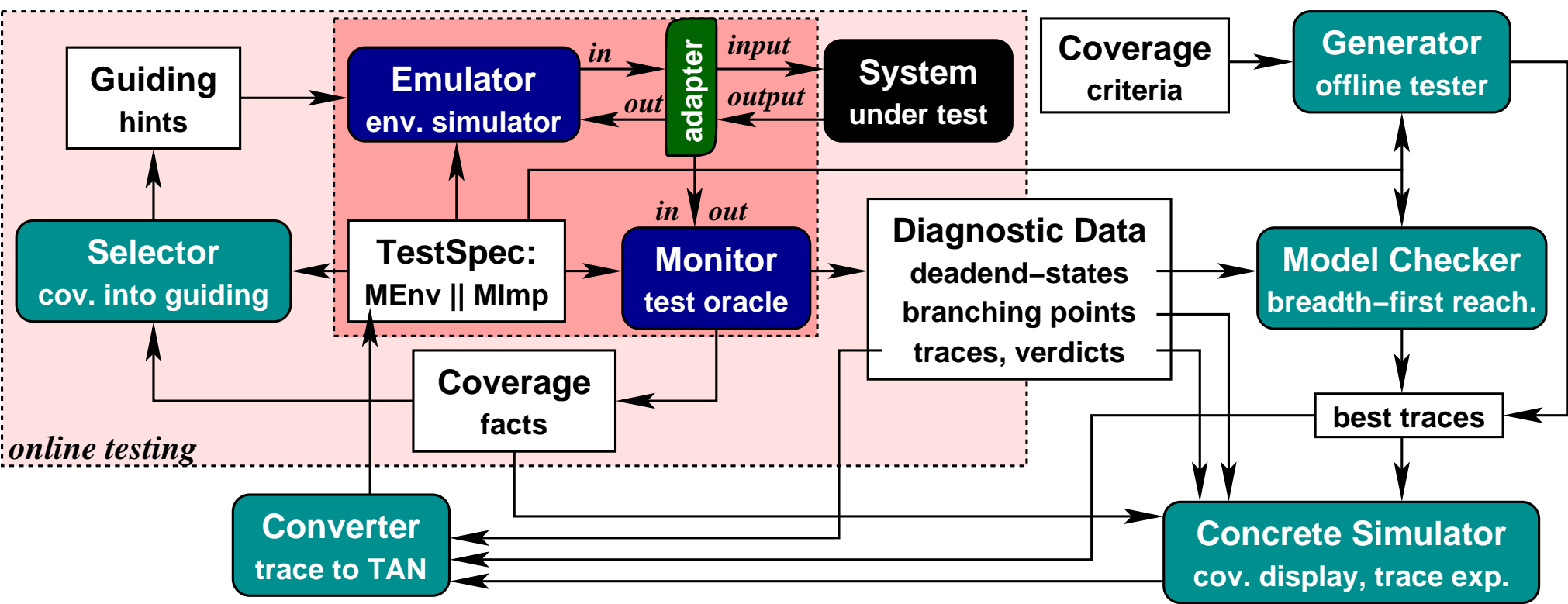
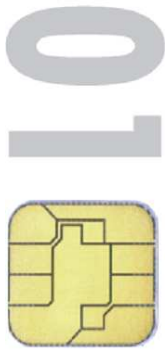
Summary and Future Work



Summary and Future Work

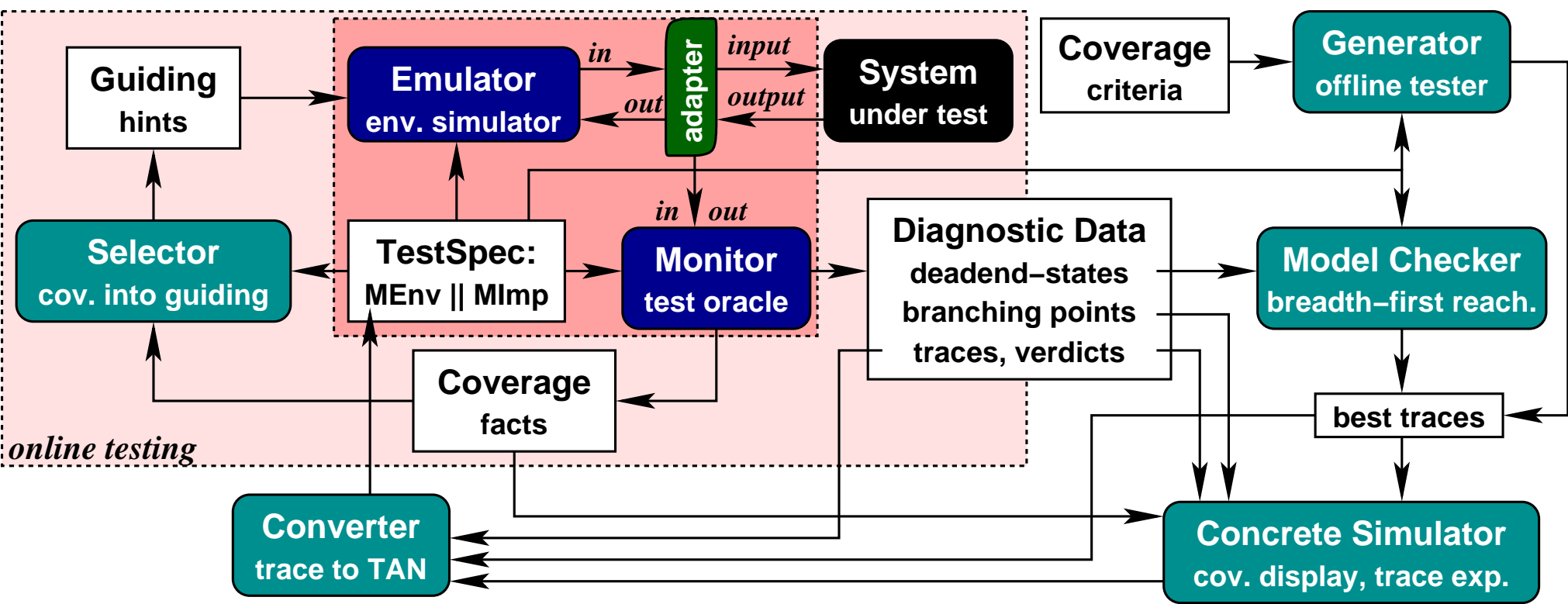
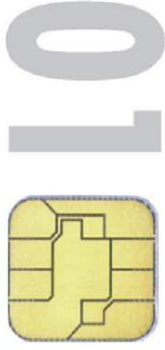


Summary and Future Work

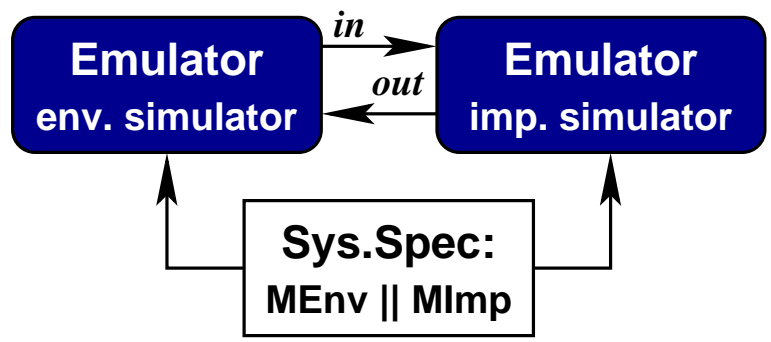


online testing

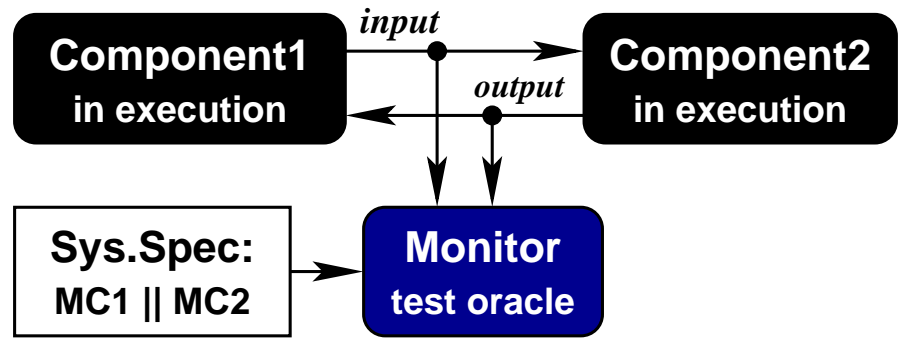
Summary and Future Work



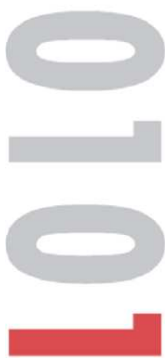
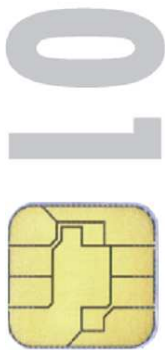
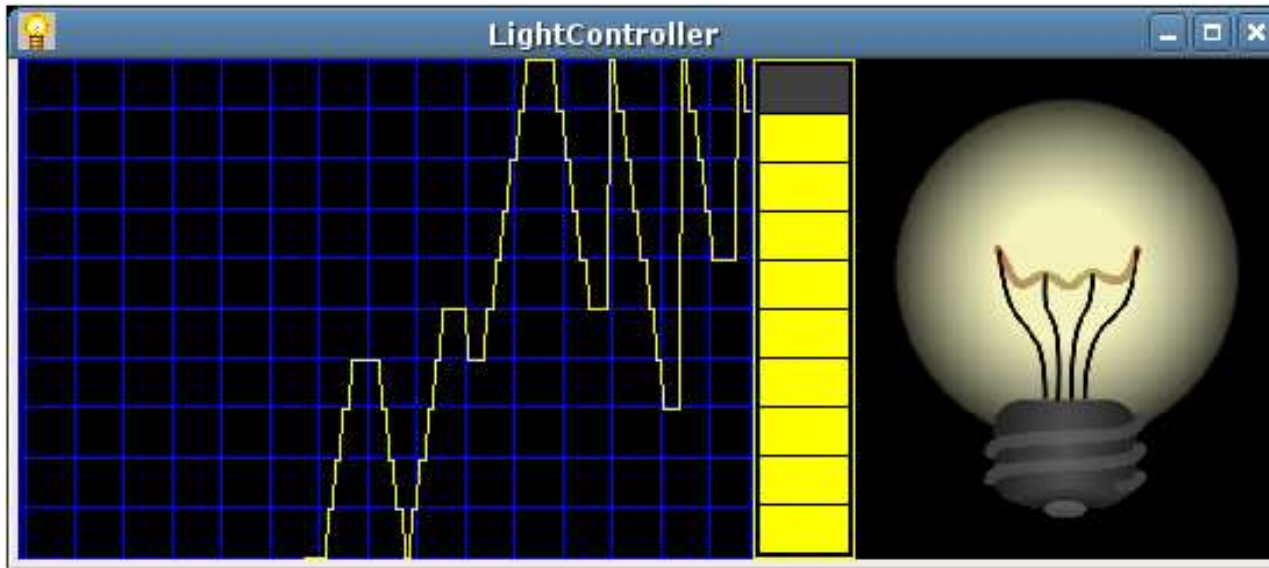
Prototyping via simulation:



Execution monitoring:



Demo of Java “Light Controller”



Supported Platforms

Compiled for the following operating systems:

- POSIX: Linux on Intel and Sun Solaris on Sparc
- Windows 2000/XP (and probably will work on others)

Currently available adapters:

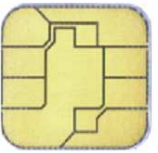
- C dynamically linked library
- Text via standard input-output streams
- TCP/IP client-server sockets:
C/C++ and Java implementations



Danfoss Refrigeration Controller



01



01



01

01



Beyond timed automata...

- From EKC manual: controller calculates the current temperature about every second via: $T_n = 80\% \cdot T_{n-1} + 20\% \cdot T_{sensor}$
- New manual: sensor values are calibrated during deployment to increase the precision from 0.5° to 0.1° using new *Pt* sensors.
- Display temperature can be roughly described by:
$$\frac{dT}{dt} = 0.2 \cdot (T_s - T) \quad \Rightarrow \quad T(t) = T_0 + \alpha_{T_s} \cdot e^{-\beta_{T_s} \cdot t}$$
- The display temperature update interval varies and sampling is performed on a soft-real-time system \Rightarrow difficult to model accurately (in timed automata with integers).
- PHAVer is more natural in modeling the dynamic behavior:
$$T_{sensor} > T : \quad 0.15 \cdot (T_{sensor} - T) \leq T' \leq 0.6 \cdot (T_{sensor} - T)$$

$$T_{sensor} < T : \quad 0.6 \cdot (T_{sensor} - T) \leq T' \leq 0.15 \cdot (T_{sensor} - T)$$



Monitoring EKC as HS using PHAVer

