

Model Based Testing

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FSM based testing

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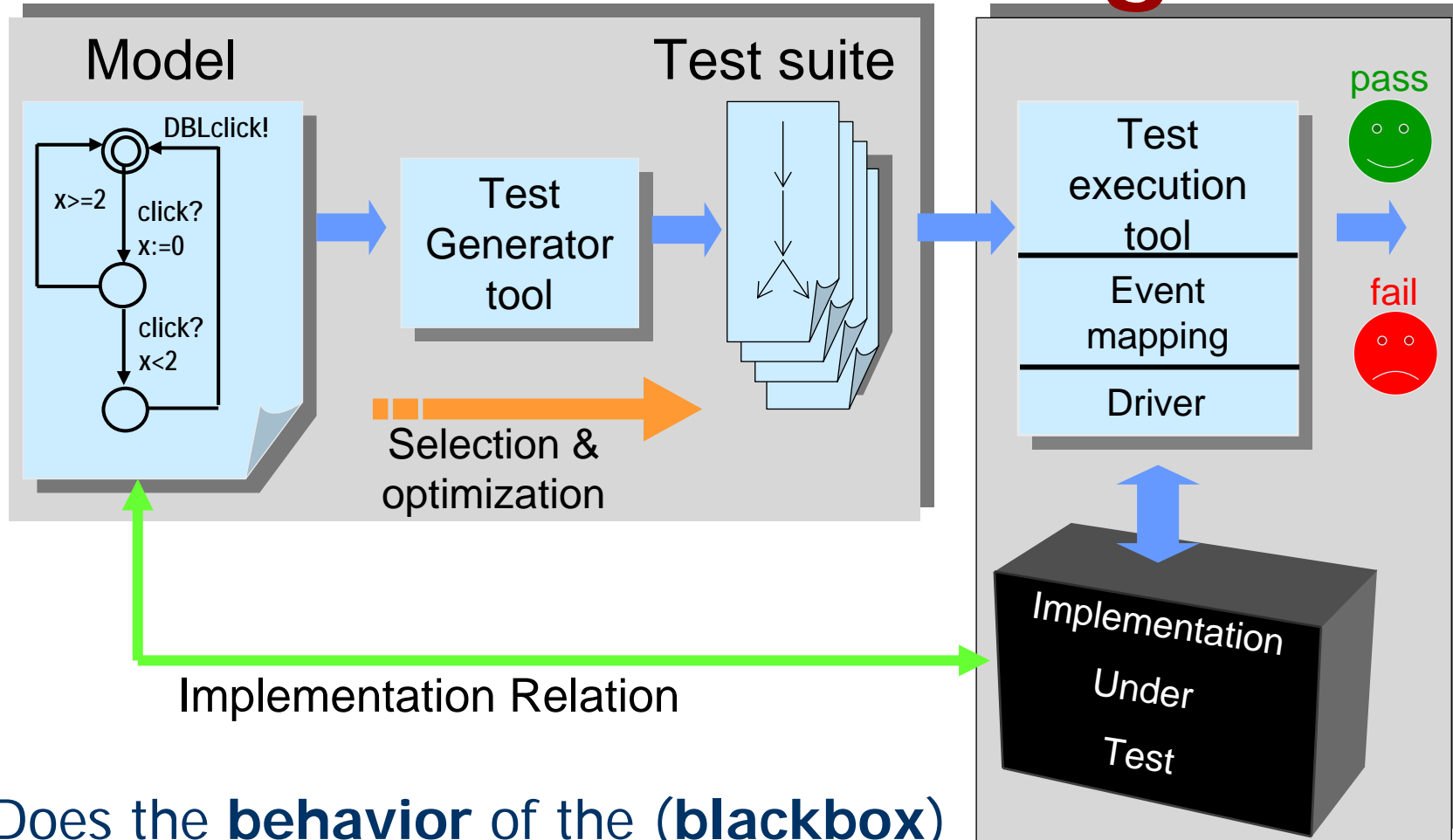
BRICS

Basic Research
in Computer Science



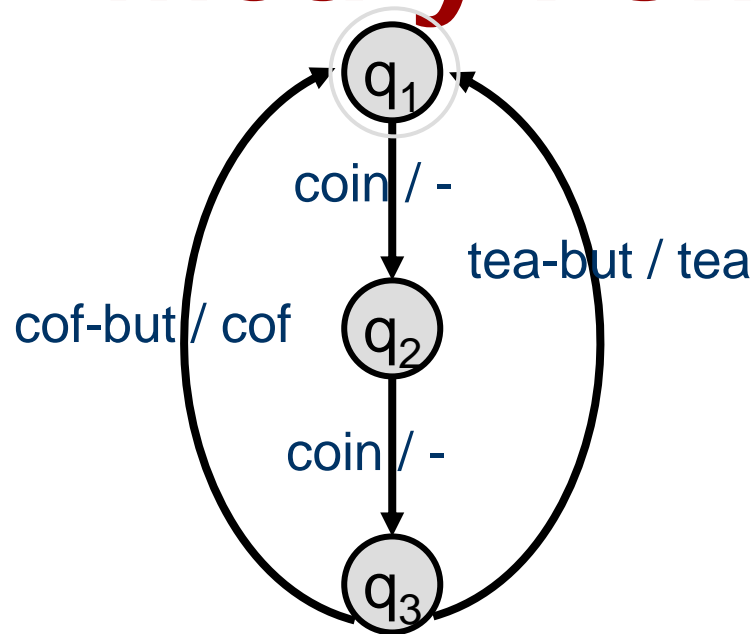
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Automated Model Based Conformance Testing



Does the **behavior** of the (**blackbox**) implementation *comply* to that of the specification?

Mealy FSM



condition		effect	
current state	input	output	next state
q ₁	coin	-	q ₂
q ₂	coin	-	q ₃
q ₃	cof-but	cof	q ₁
q ₃	tea-but	tea	q ₁

Inputs = {cof-but, tea-but, coin}

Outputs = {cof, tea}

States: {q₁, q₂, q₃}

Initial state = q₁

Transitions= {

(q₁, coin, -, q₂),

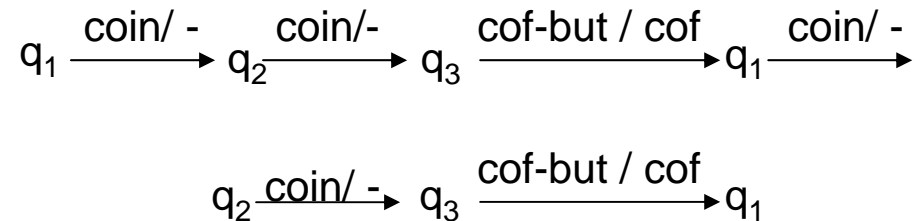
(q₂, coin, -, q₃),

(q₃, cof-but, cof, q₁),

(q₃, tea-but, tea, q₁)

}

Sample run:



State Machine : FSM Model

FSM - Finite State Machine - or *Mealy Machine* is 5-tuple

$$M = (S, I, O, \delta, \lambda)$$

S	finite set of states
I	finite set of inputs
O	finite set of outputs
$\delta : S \times I \rightarrow S$	transfer function
$\lambda : S \times I \rightarrow O$	output function

Natural extension to sequences :

$$\delta : S \times I^* \rightarrow S$$
$$\lambda : S \times I^* \rightarrow O^*$$

Concepts

- Two states s and t are (language) **equivalent** iff
 - ✱ s and t accepts same language
 - ✱ has same traces: $tr(s) = tr(t)$
- Two Machines M_0 and M_1 are equivalent iff initial states are equivalent
- A **minimized** / reduced M is one that has no equivalent states
 - ✱ for no two states $s, t, s \neq t, s$ equivalent t

Fundamental Results

- Every FSM may be determinized accepting the same language (potential explosion in size).
- For each FSM there exist a language-equivalent *minimal* deterministic FSM.
- FSM's are closed under \cap and \cup
- FSM's may be described as regular expressions (and vice versa)

Conformance Testing



Given a specification FSM M_S

a (black box) implementation FSM M_I

determine whether M_I conforms to M_S .

i.e., M_I behaves in accordance with M_S

i.e., whether outputs of M_I are the same as of M_S

i.e., whether the reduced M_I is equivalent to M_S

Today:

- Deterministic Specifications
- SUT is an (unknown) deterministic FSM (testing hypothesis)

Restrictions

FSM restrictions:

- *deterministic*

$\delta : S \times I \rightarrow S$ and $\lambda : S \times I \rightarrow O$ are *functions*

- *completely specified*

$\delta : S \times I \rightarrow S$ and $\lambda : S \times I \rightarrow O$ are *complete* functions

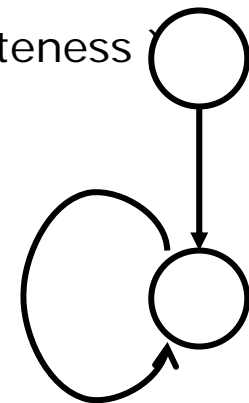
(empty output is allowed; sometimes implicit completeness)

- *strongly connected*

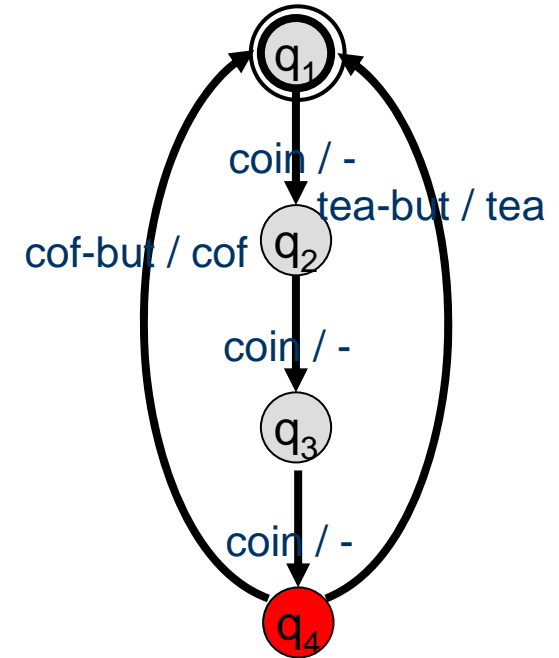
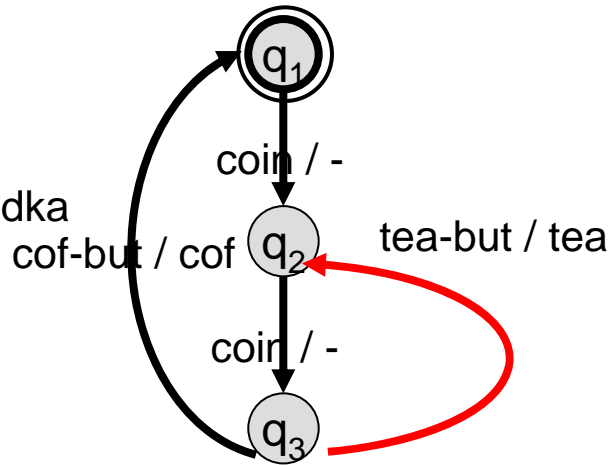
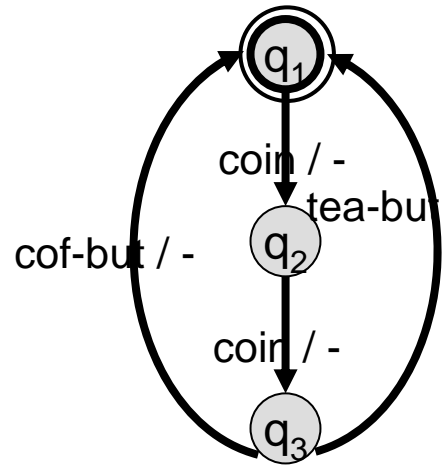
from any state any other state can be reached

- *reduced*

there are no equivalent states



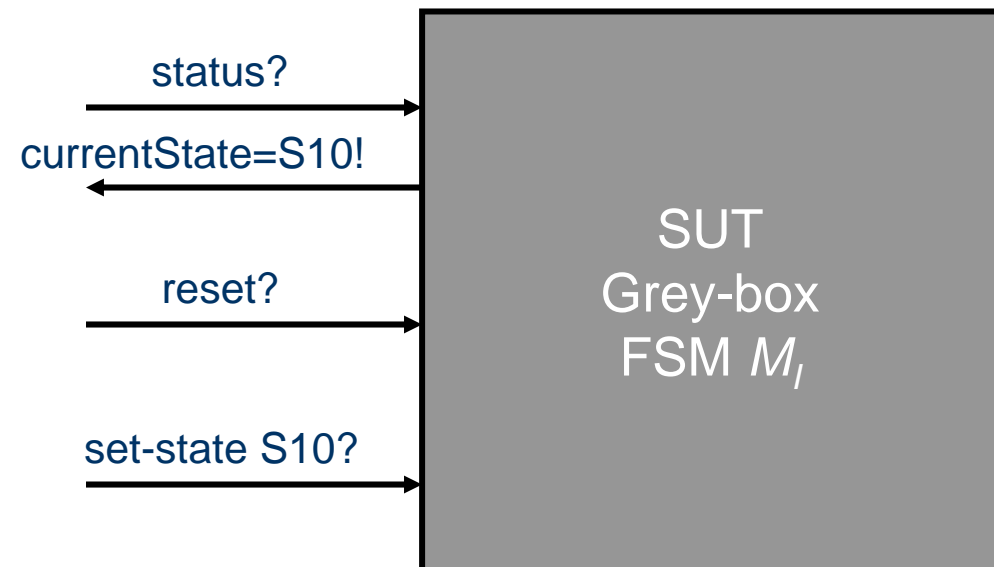
Possible Errors



- output fault (wrong or missing)
- extra or missing states
- transition fault
 - to other state
 - to new state

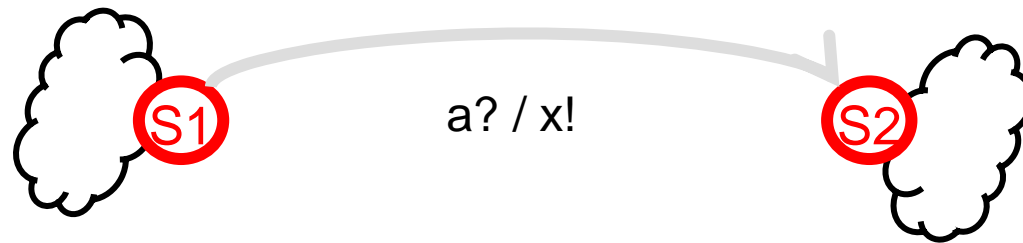
Desired Properties

- Nice, but rare / problematic
 - ✿ status messages: Assume that tester can ask implementation for its current state (reliably!!) without changing state
 - ✿ reset: reliably bring SUT to initial state
 - ✿ set-state: reliably bring SUT to any given state



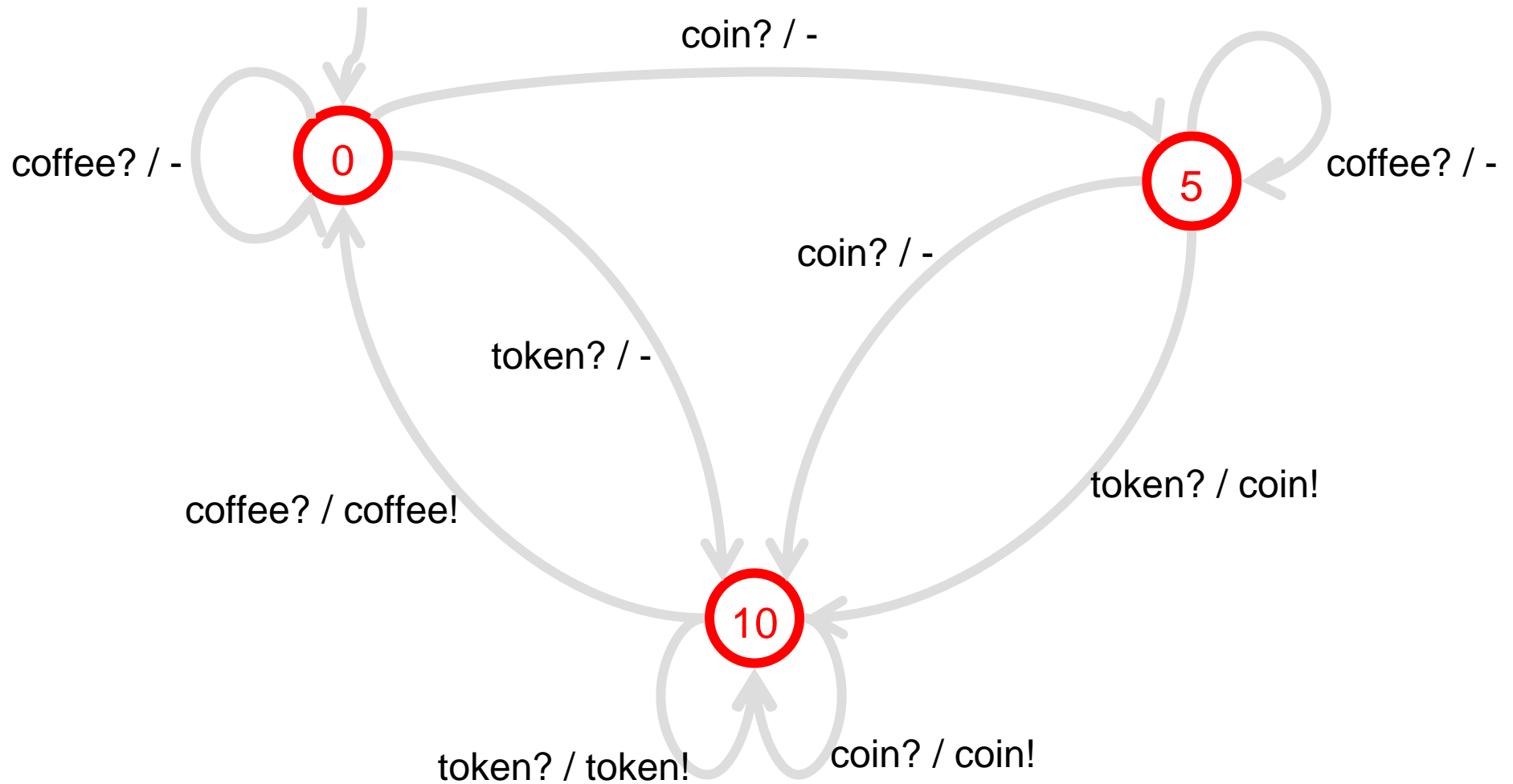
FSM Transition Testing

- Make test case for every transition in spec separately:



- Test transition :
 1. Go to state S1
 2. Apply input a?
 3. Check output x!
 4. Verify state S2 (optionally)
- *Test purpose: "Test whether the system, when in state S1, produces output x! on input a? and goes to state S2"*

Coffee Machine FSM Model



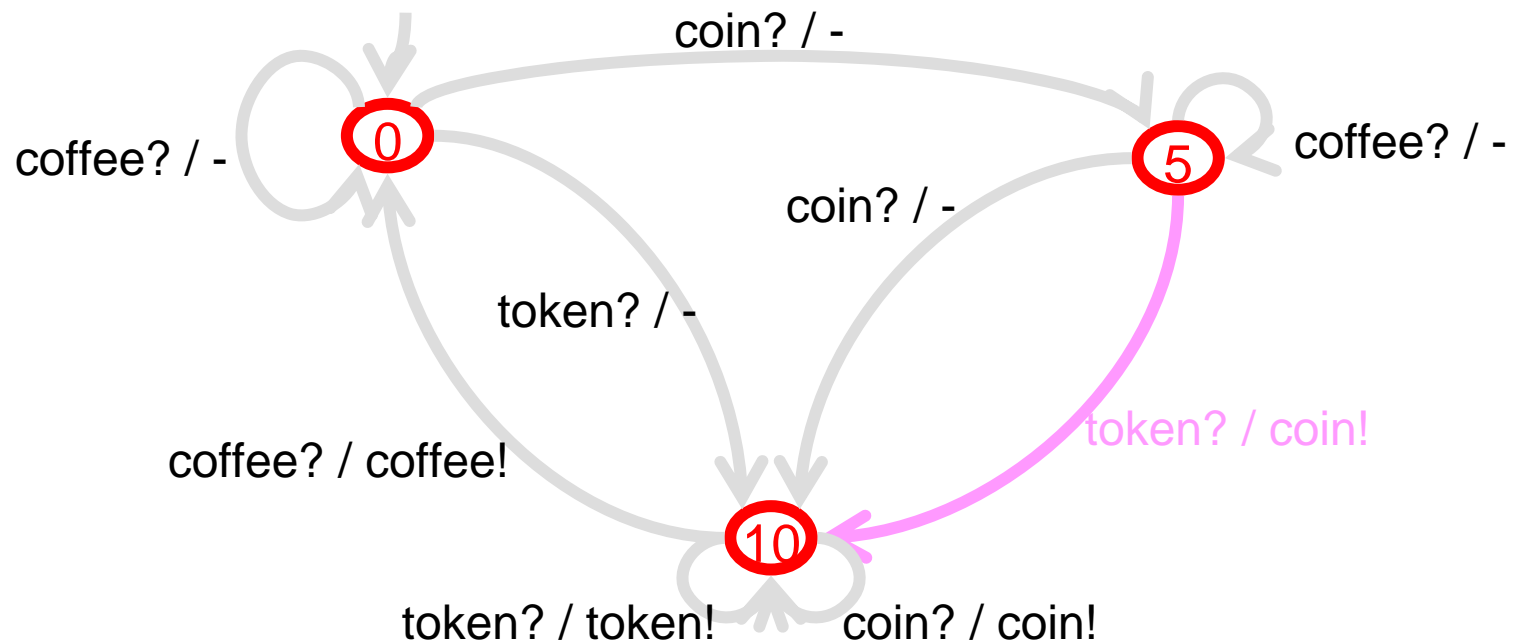
Transition Testing –1

•To test `token? / coin!` :

go to state `5` : `set-state 5`

give input `token?` check output `coin!`

verify state: send `status?` check `status=10`

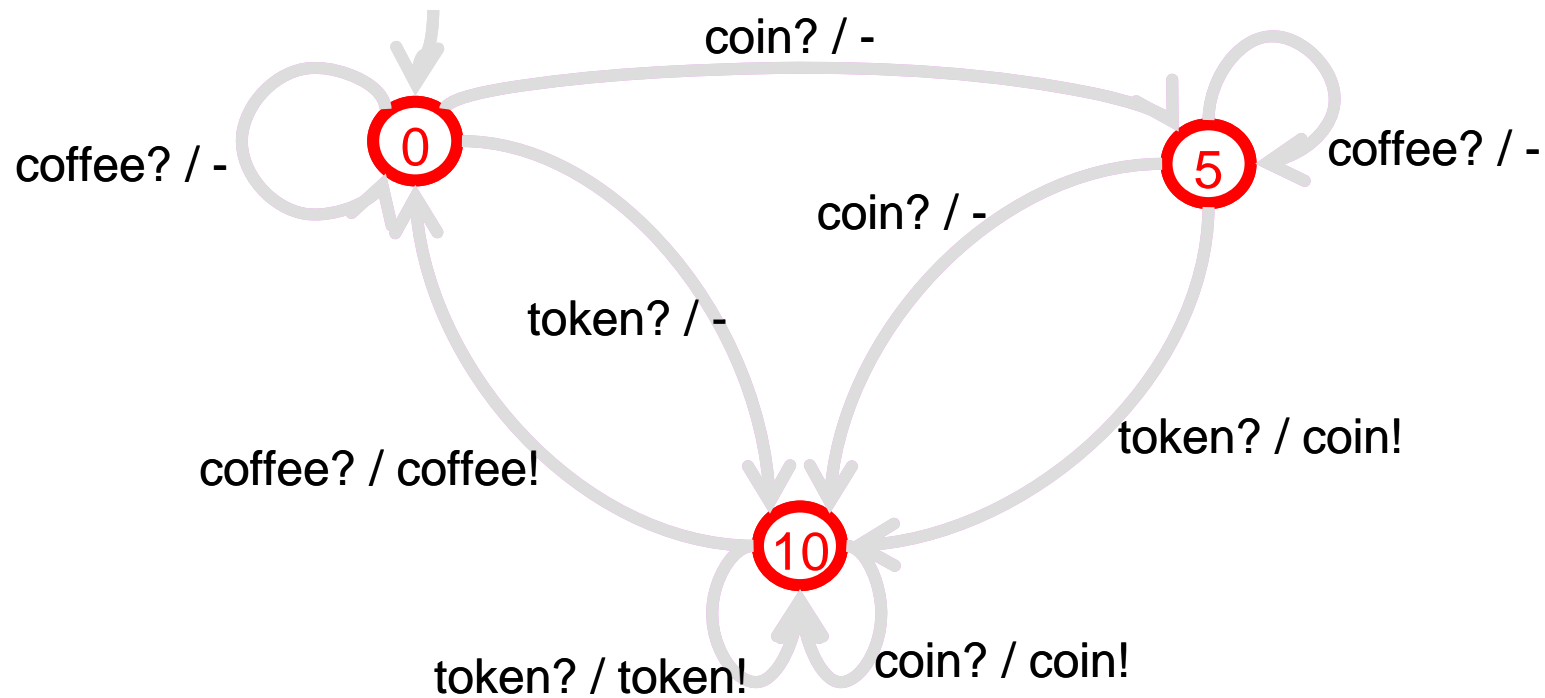


Test case : `set-state 5/ * - token? / coin! - status? / 10!`

`|S| * |I|` test cases remaining

FSM Transition Tour

- Make Transition Tour that covers every transition (in spec)



Test input sequence :

reset? coffee? coin? coffee? coin? coin? token? coffee? token? coffee? coin? token? coffee?

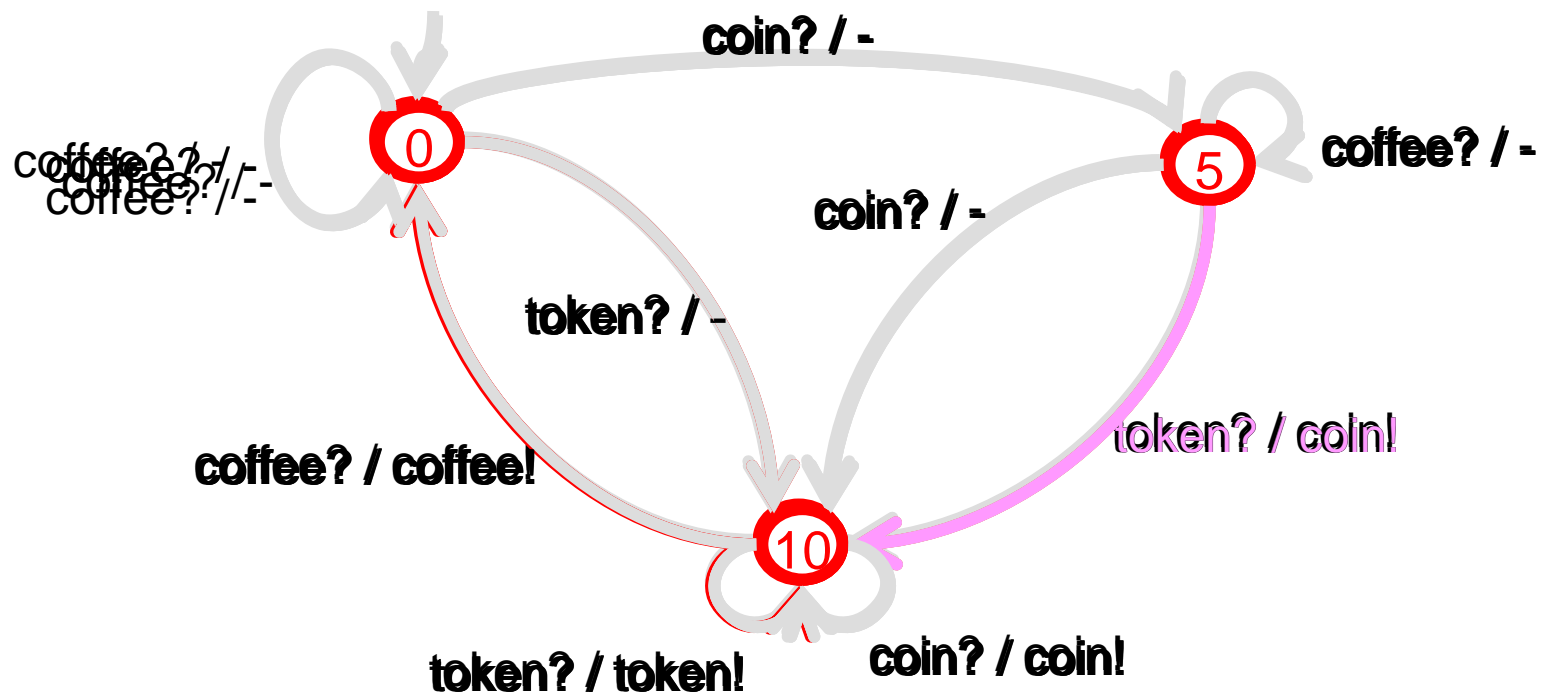
+ check **expected outputs** and target state by **status message**

Transition Testing -1

- Go to state S5 :
- No Set-state property???
 - ✱ use *reset property* if available
 - ✱ go from S0 to S5
(always possible because of determinism and completeness)
 - ✱ or:
 - ✱ *synchronizing sequence* brings machine to particular known state, say S0, from any state
 - ✱ (but synchronizing sequence may not exist)

Transition Testing -1

synchronizing sequence : token? coffee?

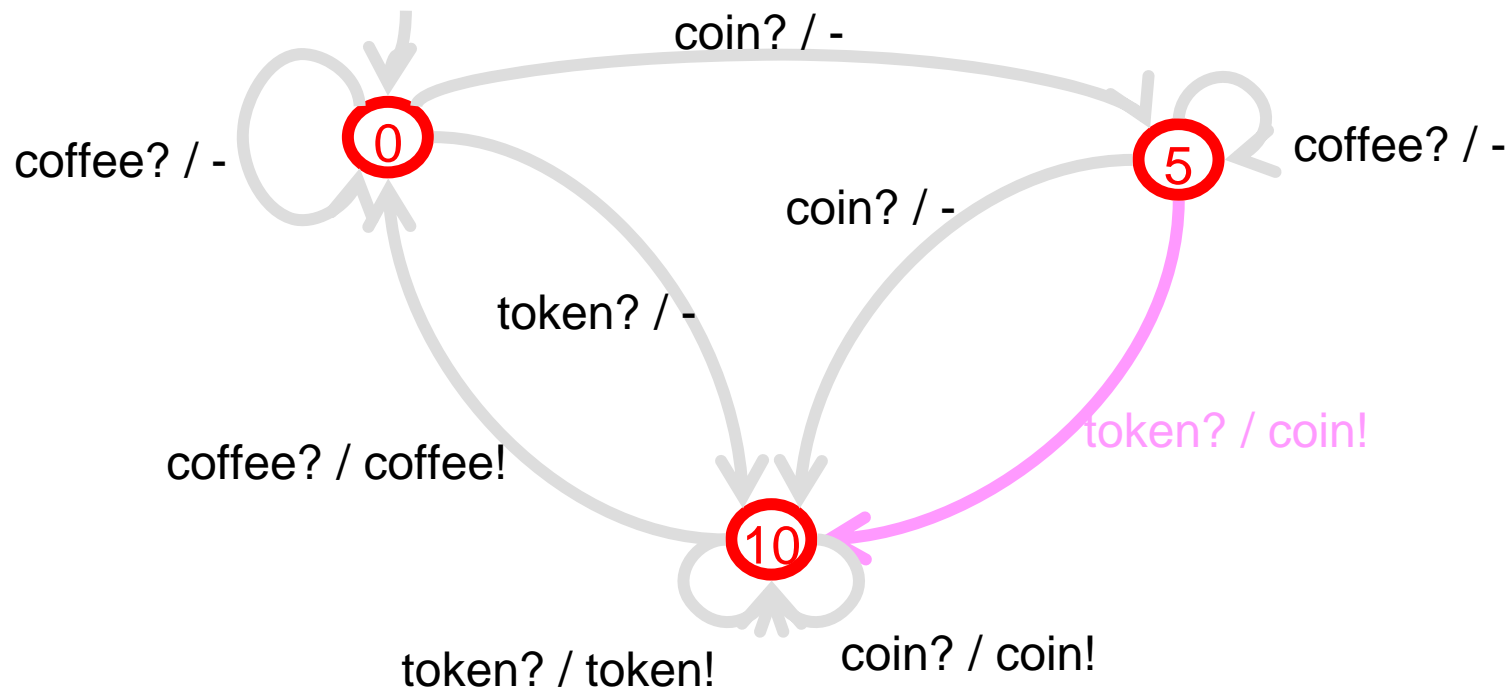


To test **token? / coin!** : go to state **5** by : token? coffee? coin?

Transition Testing –2,3

•To test *token? / coin!* :

1. go to state **5** by : *token? coffee? coin?*
2. give input *token?*
3. check output *coin!*
4. verify that machine is in state **10**



Transition Testing-4

- No Status Messages??
- **State identification: What state am I in??**
- **State verification : Am I in state s?**
 - ✱ Apply sequence of inputs in the current state of the FSM such that from the outputs we can
 - identify that state where we started; or
 - verify that we were in a particular start state
 - ✱ Different kinds of sequences
 - UIO sequences (Unique Input Output sequence, SIOS)
 - Distinguishing sequence (DS)
 - W - set (characterizing set of sequences)
 - UIOv
 - SUIO
 - MUIO
 - Overlapping UIO

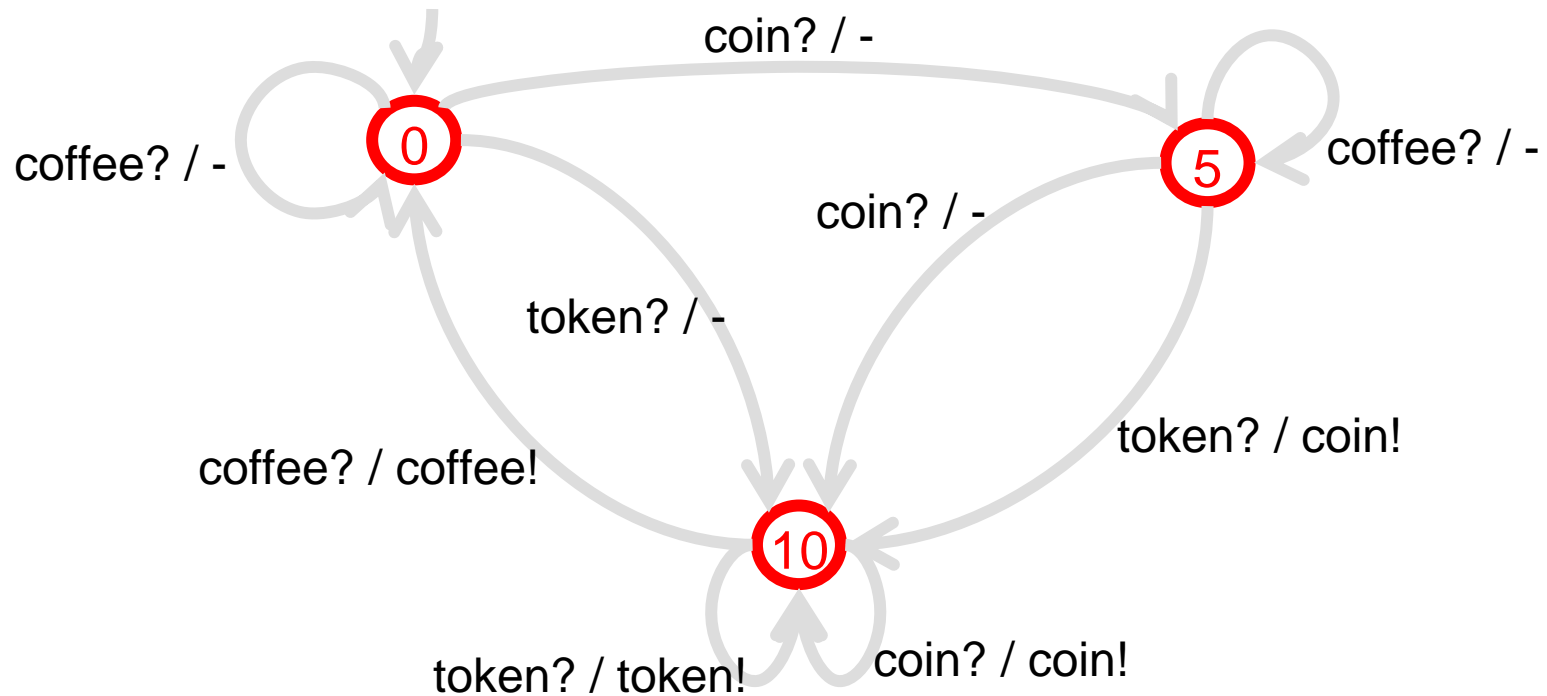
Transition Testing-4

State check :

- UIO sequences (verification)
 - ✱ sequence x_s that distinguishes state s from all other states :
for all $t \neq s$: $\lambda (s, x_s) \neq \lambda (t, x_s)$
 - ✱ each state has its own UIO sequence
 - ✱ UIO sequences may not exist
- Distinguishing sequence (identification)
 - ✱ sequence x that produces different output for every state :
for all pairs t, s with $t \neq s$: $\lambda (s, x) \neq \lambda (t, x)$
 - ✱ a distinguishing sequence may not exist
- W - set of sequences (identification)
 - ✱ *set of* sequences W which can distinguish any pair of states :
for all pairs $t \neq s$ there is $x \in W$: $\lambda (s, x) \neq \lambda (t, x)$
 - ✱ W - set always exists for reduced FSM

Transition Testing-4: UIO

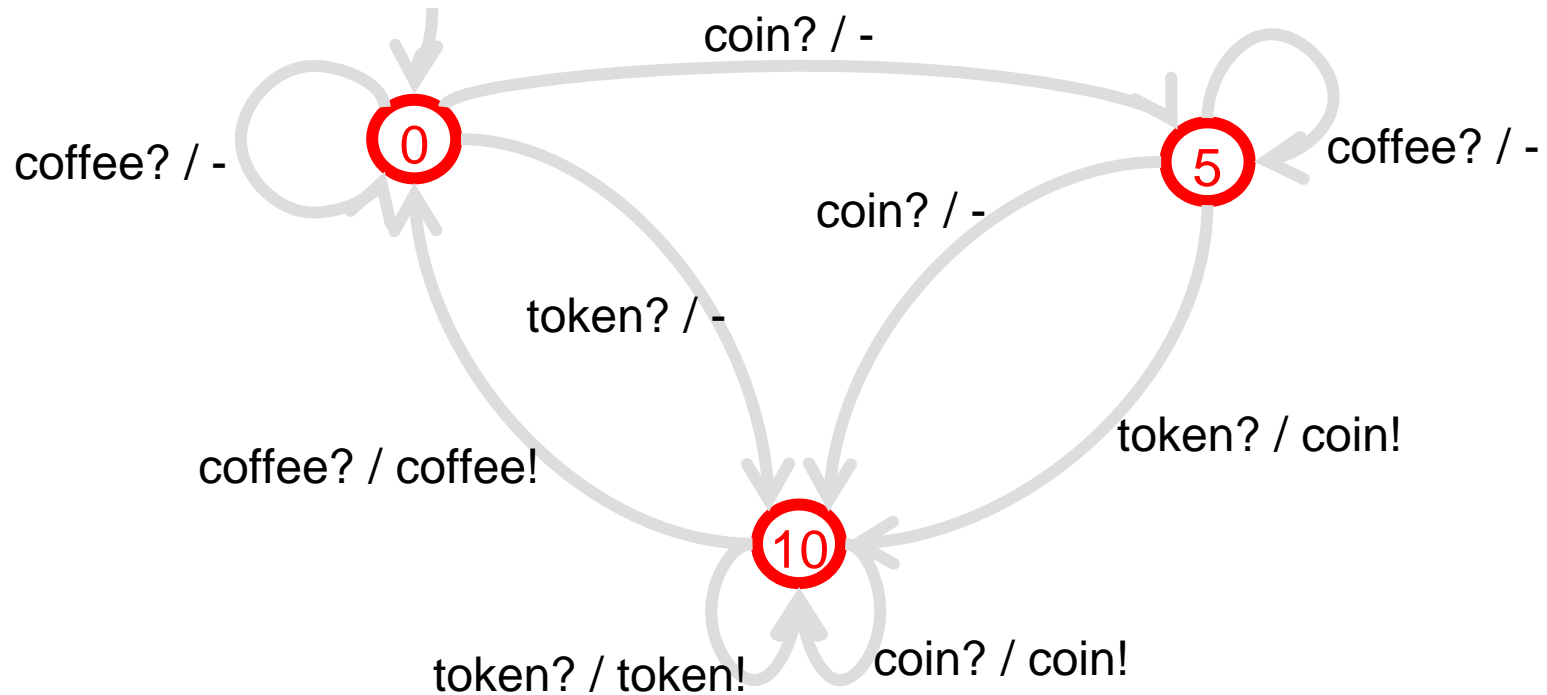
UIO sequences



state 0 : coin? / - coffee? / -
state 5 : token? / coin!
state 10 : coffee? / coffee!

Transition Testing-4: DS

DS sequence



DS sequence : token? output state 0 : -
output state 5 : coin!
output state 10 : token!

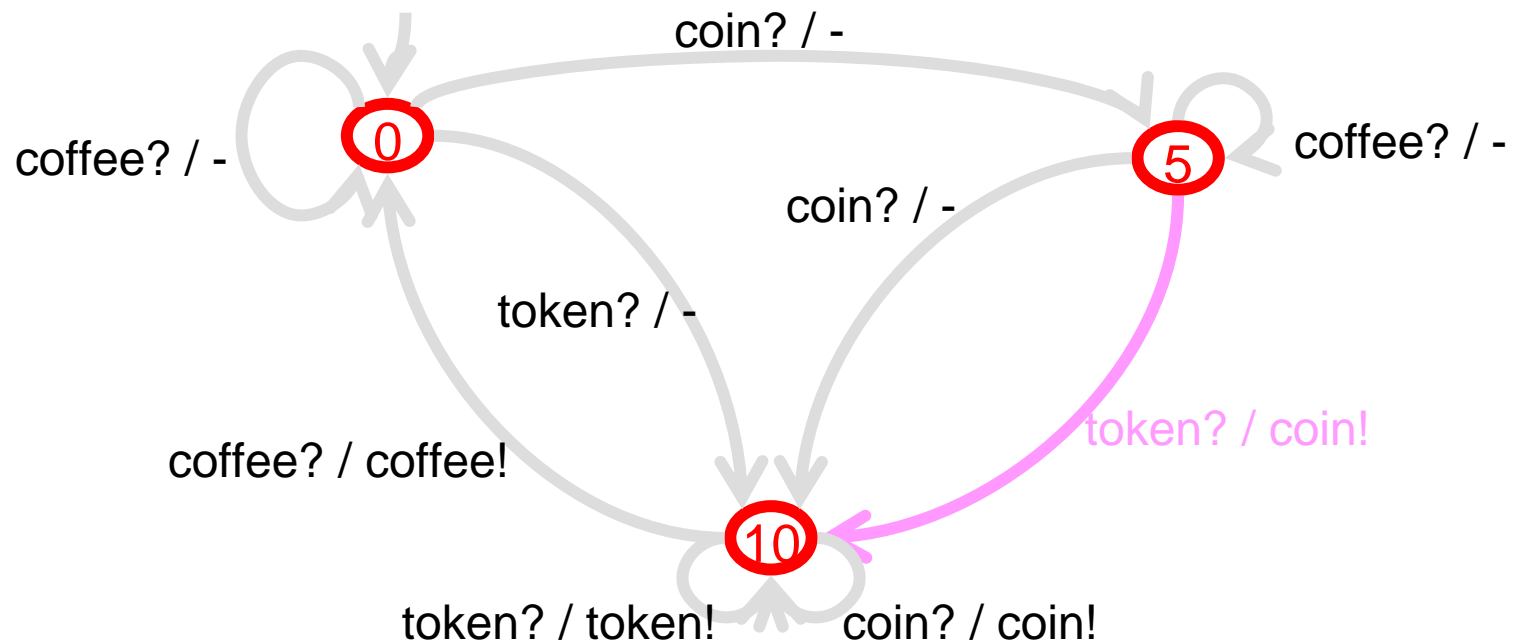
Transition Testing –4 done

- To test **token? / coin!** :

go to state **5** : token? coffee? coin?

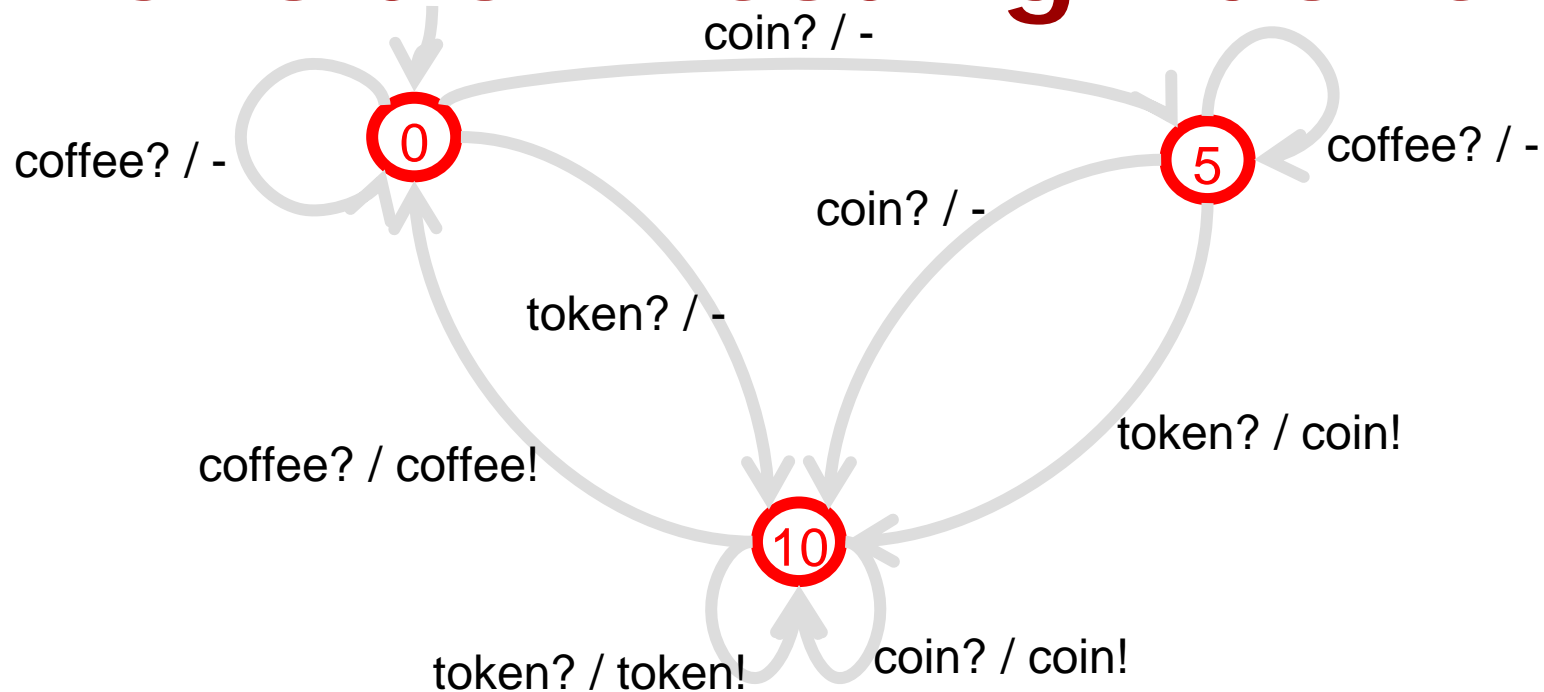
give input **token?** check output **coin!**

Apply UIO of state **10** : coffee? / coffee!



Test case : token? / * coffee? / * coin? / - **token? / coin!** coffee? / coffee!

Transition Testing - done



- 9 transitions / test cases for coffee machine
- if end-state of one corresponds with start-state of next then concatenate
- different ways to optimize and remove overlapping / redundant parts
- there are (academic) tools to support this

FSM Transition Testing

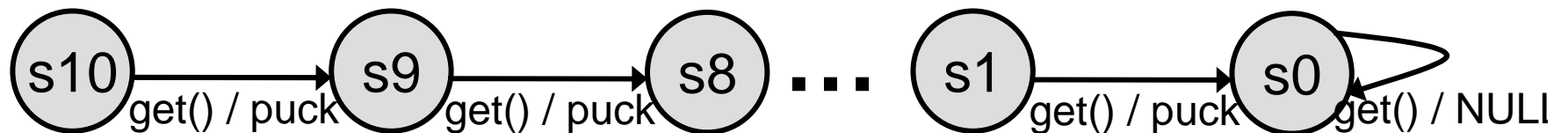
- Test transition :
 - Go to state S_1
 - Apply input a ?
 - Check output $x!$
 - Verify state S_2
- Checks every output fault and transfer fault (to existing state)
- **If** we assume that
 - the number of states of the implementation machine M_i*
 - is less than or equal to*
 - the number of states of the specification machine to M_s .*then testing all transitions in this way leads to equivalence of reduced machines, i.e., **complete conformance**
- If not: exponential growth in test length in number of extra states.

Object Testing-1

```
Class PuckSupply{  
  int _count=10;  
Public:  
  
  puck * get(){  
    if(_count>0) {  
      _count--;return get_puck();  
    } else {  
      return NULL;  
    }  
  };  
};
```

- A Test case

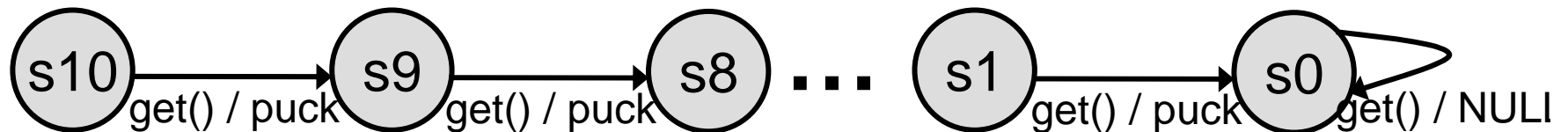
```
s=new PuckSupply;  
Assert(s.get()== aPuck);  
Assert(s.get()== aPuck);  
Assert(s.get()== aPuck);  
Assert(s.get()== aPuck);  
Assert(s.get()== aPuck);  
Assert(s.get()== aPuck);  
Assert(s.get()== aPuck);  
Assert(s.get()== aPuck);  
Assert(s.get()== aPuck);  
Assert(s.get()== aPuck);  
Assert(s.get()== aPuck);  
Assert(s.get()== aPuck);  
Assert(s.get()== aPuck);  
Assert(s.get()== NULL);  
Assert(s.get()== NULL);
```



Object Testing-2

```
Class PuckSupply{
  int _count=10;
Public:
  //test helpers
  int getState() {return _count;}
  void setState(int count) {
    _count=count;
    alloc_puks();
  }
  puck * get(){
    if(_count>0) {
      _count--;return new puck();
    } else {
      return NULL;
    }
  }
};
```

```
Assert(s.get()==aPuck);
Assert(s.getState()==9);
s.setState(1);
Assert(s.get()==aPuck);
Assert(s.getState()==0);
Assert(s.get()==NULL);
Assert(s.getState()==0);
Assert(s.get()==NULL);
Assert(s.getState()==0);
```

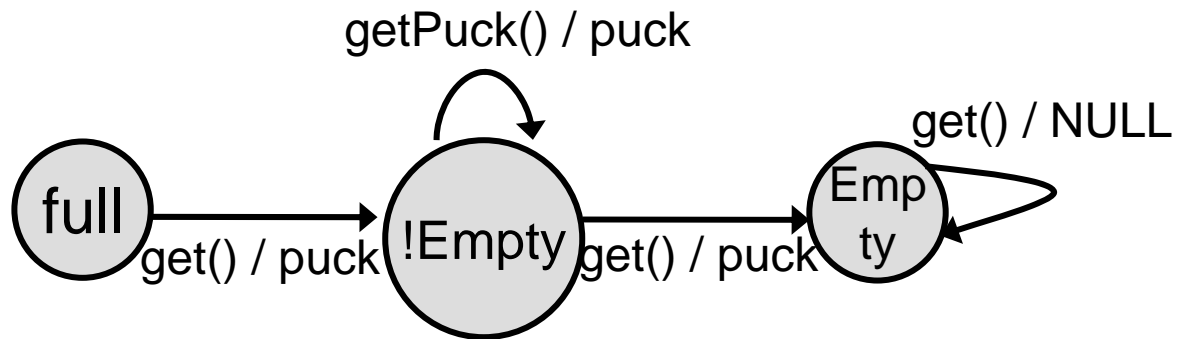


Object Testing: Abstraction

```
Class PuckSupply{
  int _count=10;
Public:

  puck * get(){
    if(_count>0) {
      _count--;return new puck();
    } else {
      return NULL;
    }
  };
};
```

How many states in corresponding FSM?

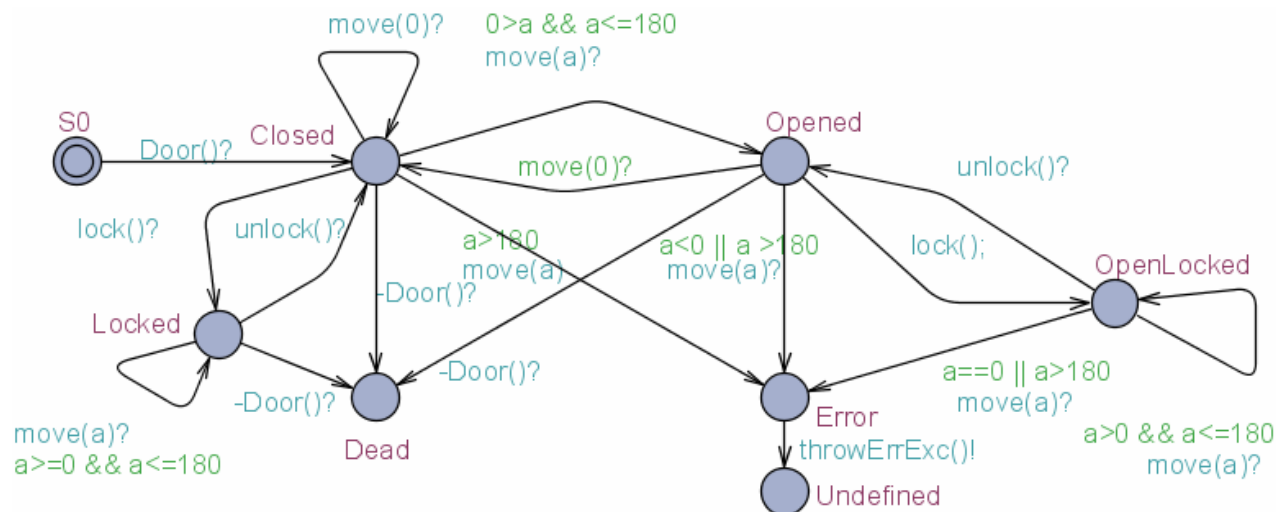


- ⇒ Generate tests systematically from **abstract** descriptions to **select** reasonably number of tests

Object Tests

- `D=new Door();`

```
Class Door{
Private:
    //state variables
    //methods
Public:
    Door();
    ~Door();
    Lock();
    Unlock();
    Move(Angle a) throws ErrorExc;
//test Helpers?
    State getState();
    void setState(State);
    void reset();}
```

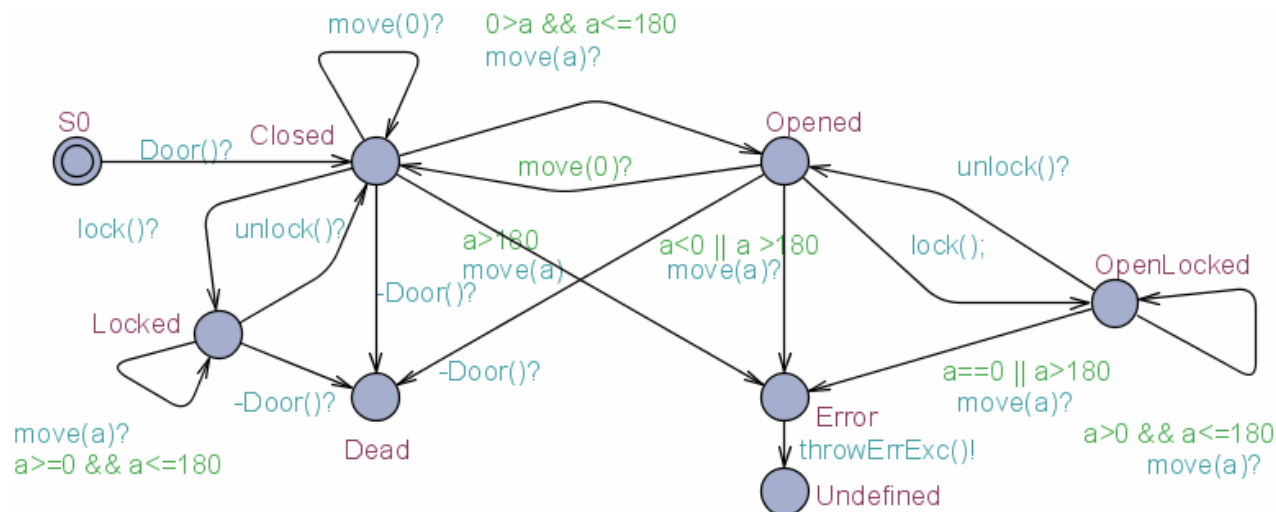


Object Tests

Test Purpose: A specific test objective (or observation) the tester wants to make on SUT

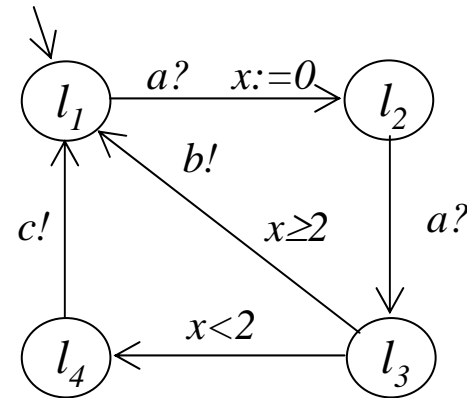
TP1: check that door can be open and locked?

- $E \leftrightarrow \text{door.OpenLocked}$
- Shortest Trace: `Door()?.move(1)?.lock()?`



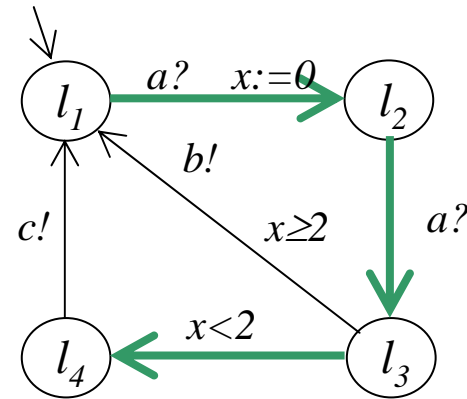
Coverage Based Test Generation

- Multi purpose testing
- Cover measurement
- Examples:
 - ✱ Location coverage,
 - ✱ Edge coverage,
 - ✱ Definition/use pair coverage



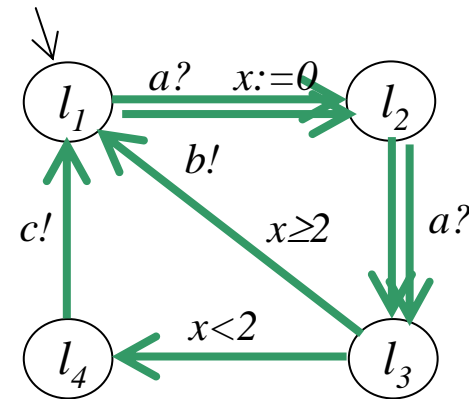
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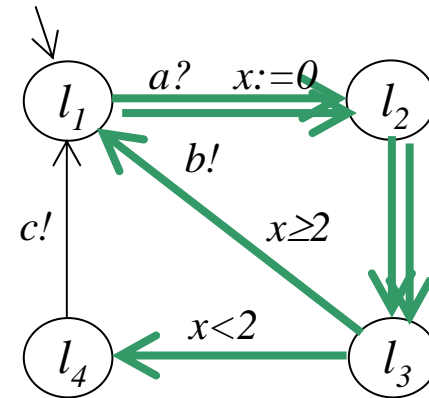
Coverage Based Test Generation

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- Examples:

- ✱ Location Coverage,

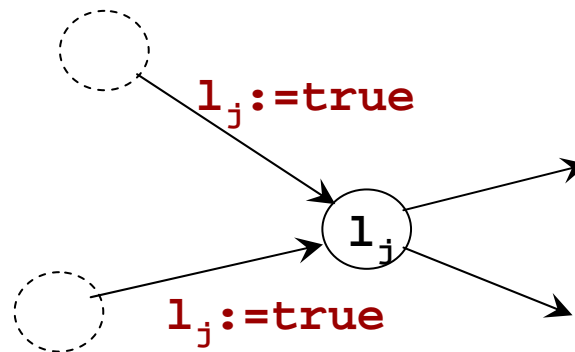
- ✱ Edge Coverage,

- ✱ **Definition/Use Pair Coverage**



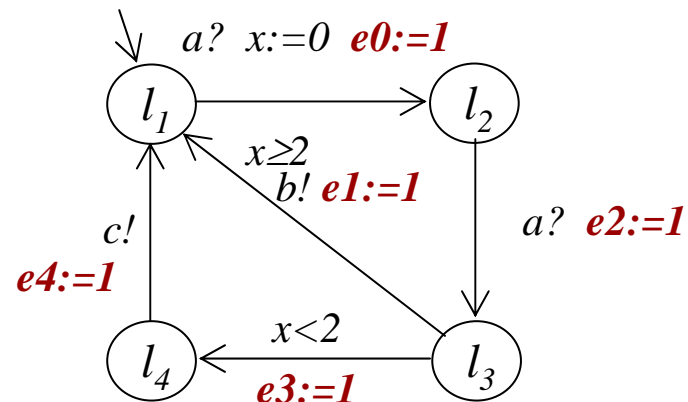
Location Coverage

- Test sequence traversing all locations
- Encoding:
 - ✱ Enumerate locations l_0, \dots, l_n
 - ✱ Add an auxiliary variable l_i for each location
 - ✱ Label each ingoing edge to location i $l_i := \text{true}$
 - ✱ Mark initial visited $l_0 := \text{true}$
- Check: $E \langle \rangle (l_0 = \text{true} \wedge \dots \wedge l_n = \text{true})$



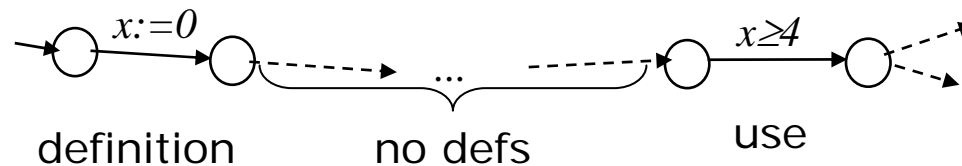
Edge Coverage

- Test sequence traversing all edges
- Encoding:
 - ✱ Enumerate edges e_0, \dots, e_n
 - ✱ Add auxiliary variable e_i for each edge
 - ✱ Label each edge $e_i := \text{true}$
- Check: $\mathbf{E} \langle \rangle (e_0 = \text{true} \wedge \dots \wedge e_n = \text{true})$



Definition/Use Pair Coverage

- Dataflow coverage technique
- Def/use pair of variable x :

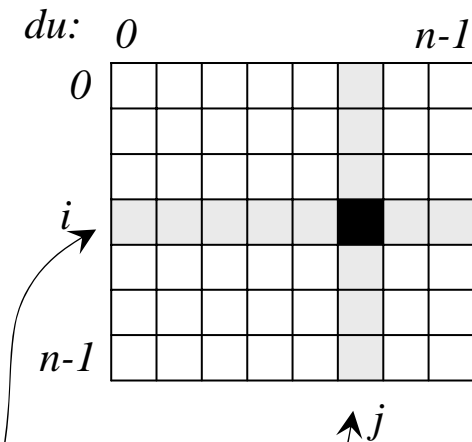
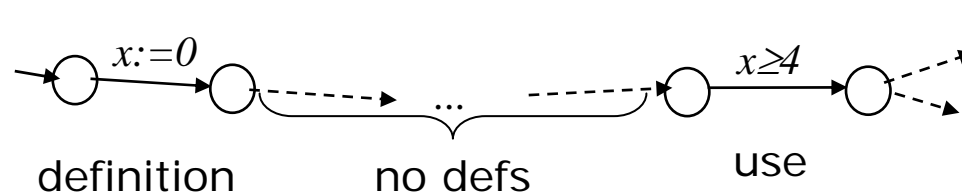


- Encoding:

- ✱ $v_d \in \{false\} \cup \{e_0, \dots, e_n\}$, initially false
- ✱ Boolean array du of size $|E| \times |E|$
- ✱ At definition on edge i : $v_d := e_i$
- ✱ At use on edge j : if(v_d) then $du[v_d, e_j] := true$

Definition/Use Pair Coverage

- Dataflow coverage technique
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- Encoding:

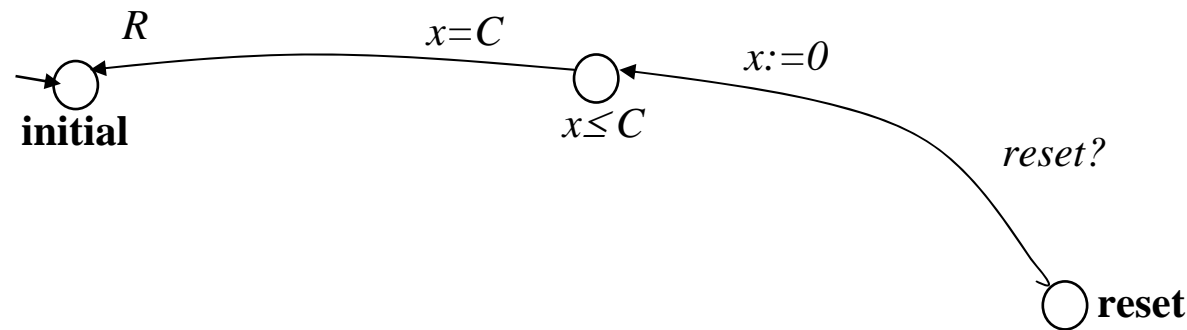
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- ✱ At definition on edge i : $v_d := e_i$
- ✱ At use on edge j : if(v_d) then $du[v_d, e_j] := true$

- Check:

- ✱ $E \ll \gg (\text{ all } du[i,j] = true)$

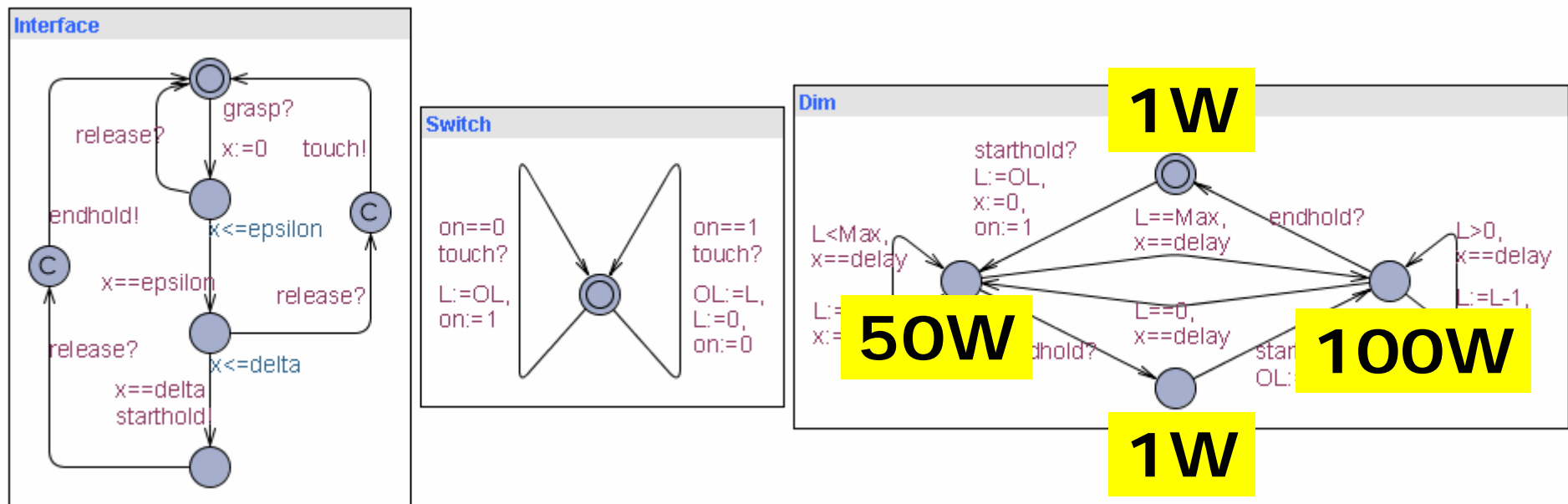
Test Suite Generation

- In general a set of test cases is needed to cover a test criteria
- Add global reset of SUT and environment model and associate a cost (of system reset)



- Same encodings and min-cost reachability
- Test sequence $\sigma = \mathcal{E}_0, i_0, \dots, \mathcal{E}_1, i_1, \mathbf{reset} \underbrace{\mathcal{E}_2, i_2, \dots, \mathcal{E}_0, i_0}_{\sigma_i}, \mathbf{reset}, \mathcal{E}_1, i_1, \mathcal{E}_2, i_2, \dots$
- Test suite $T = \{ \sigma_1, \dots, \sigma_n \}$ with minimum cost

Optimal Tests



- **Shortest** test for max light??
- **Fastest** test for max light??
- **Fastest** edge-covering test suite??
- Least **power** consuming test??